

V8 / Chromium / Chrome



Attacking JavaScript Engines in 2022

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JavaScriptCore / WebKit / Safari

<https://t.me/learningnets>

Spidermonkey / Gecko / Firefox



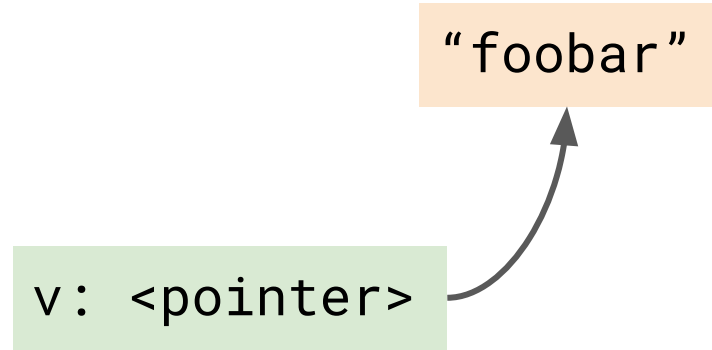
Basic JavaScript

```
let v = 0x1337;  
// typeof(v) == "number"
```

v: 0x1337

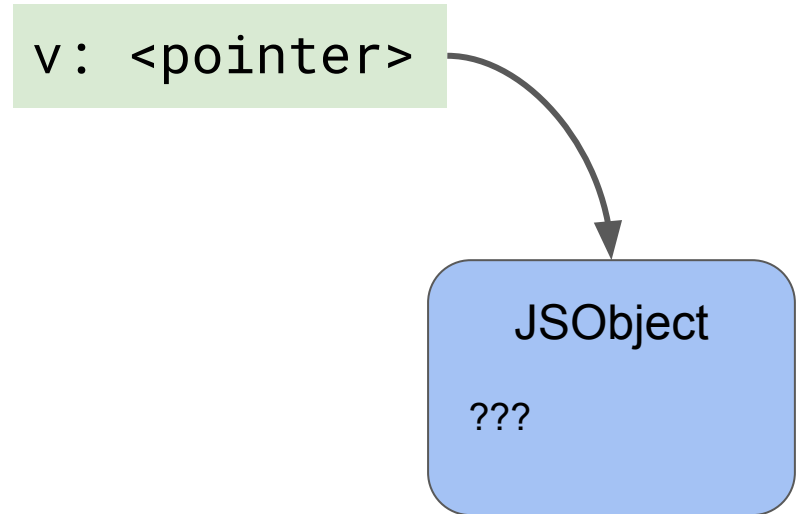
Basic JavaScript

```
let v = 0x1337;  
// typeof(v) == "number"  
v = "foobar";  
// typeof(v) == "string"
```



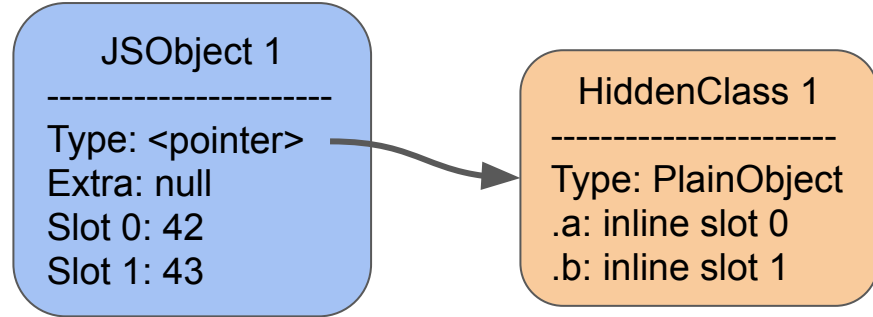
Basic JavaScript

```
let v = 0x1337;  
// typeof(v) == "number"  
v = "foobar";  
// typeof(v) == "string"  
v = {a: 42, b: 43};  
// typeof(v) == "object"
```



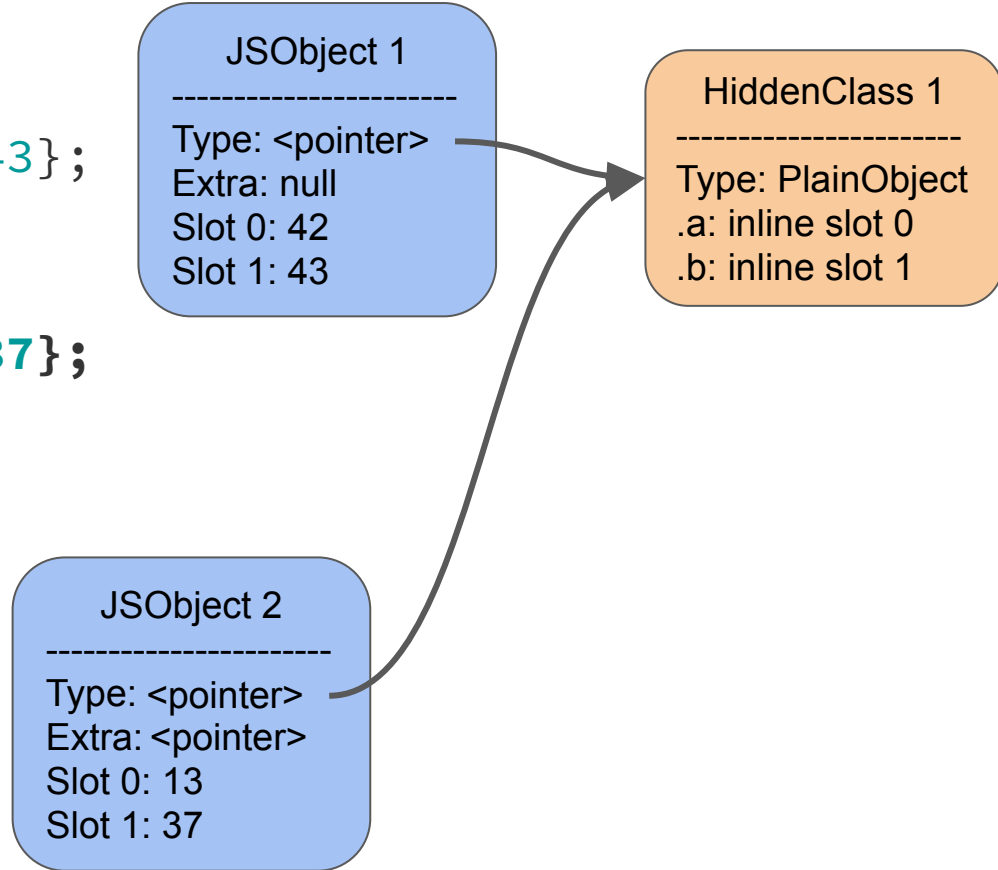
Basic JavaScript

```
let o1 = {a: 42, b: 43};  
console.log(o1.a);
```



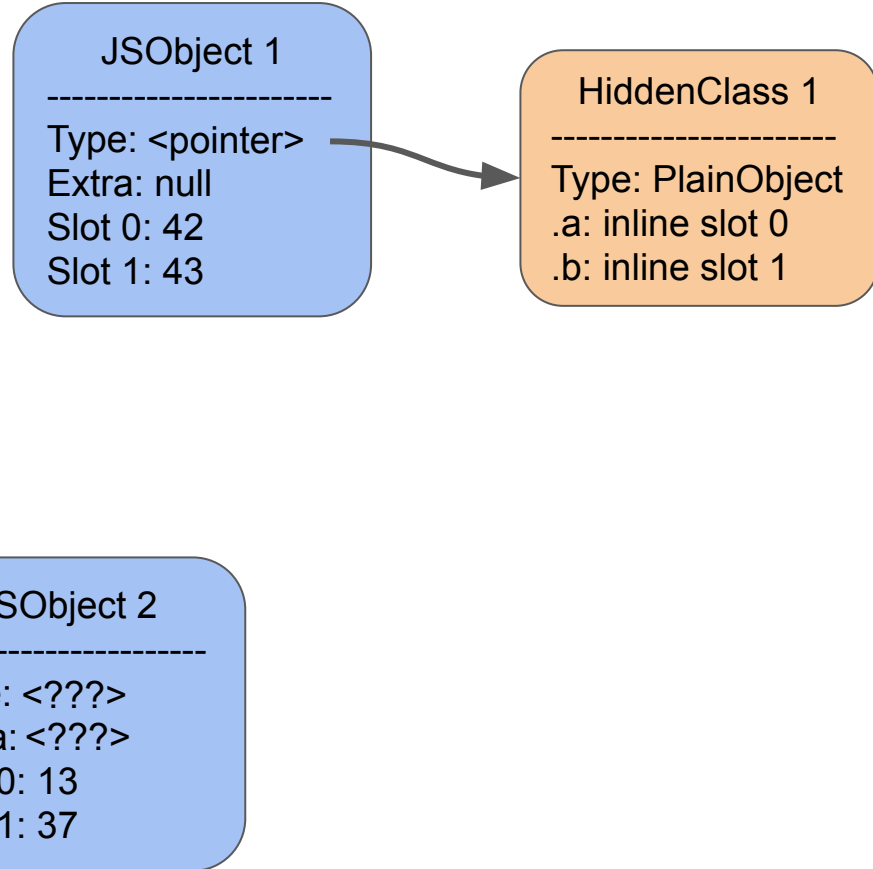
Basic JavaScript

```
let o1 = {a: 42, b: 43};  
console.log(o1.a);  
  
let o2 = {a: 13, b: 37};
```



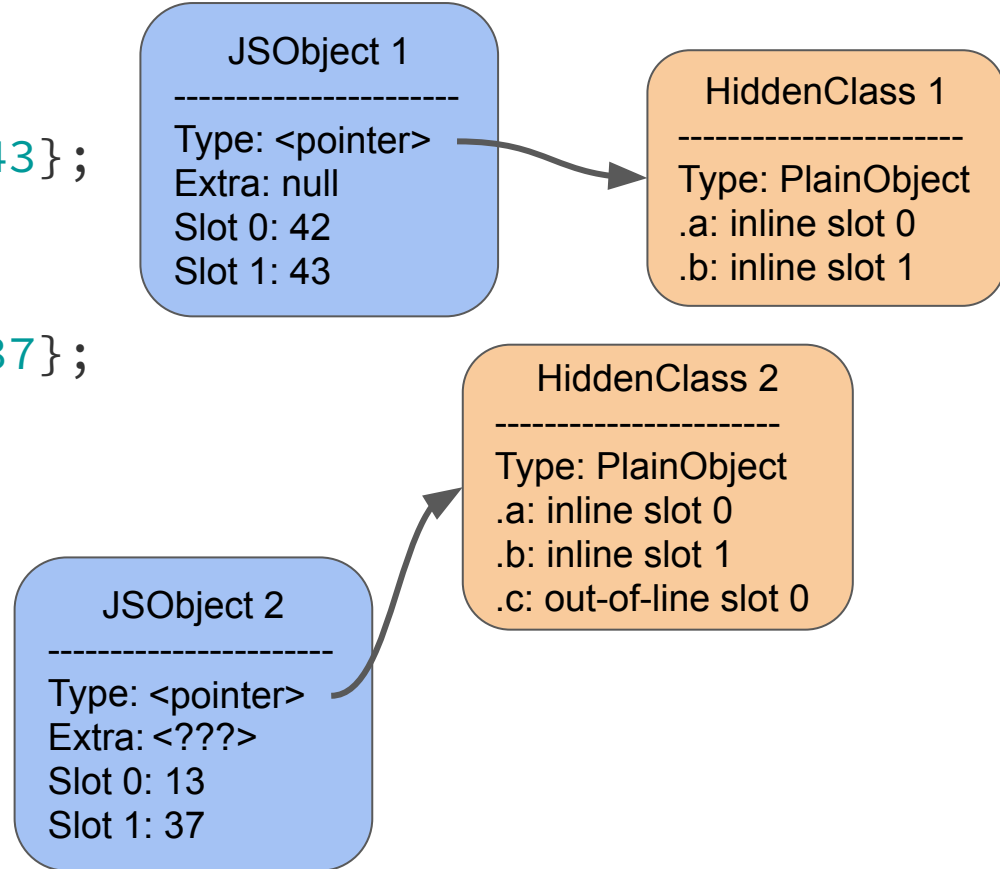
Basic JavaScript

```
let o1 = {a: 42, b: 43};  
console.log(o1.a);  
  
let o2 = {a: 13, b: 37};  
  
o2.c = o1;
```



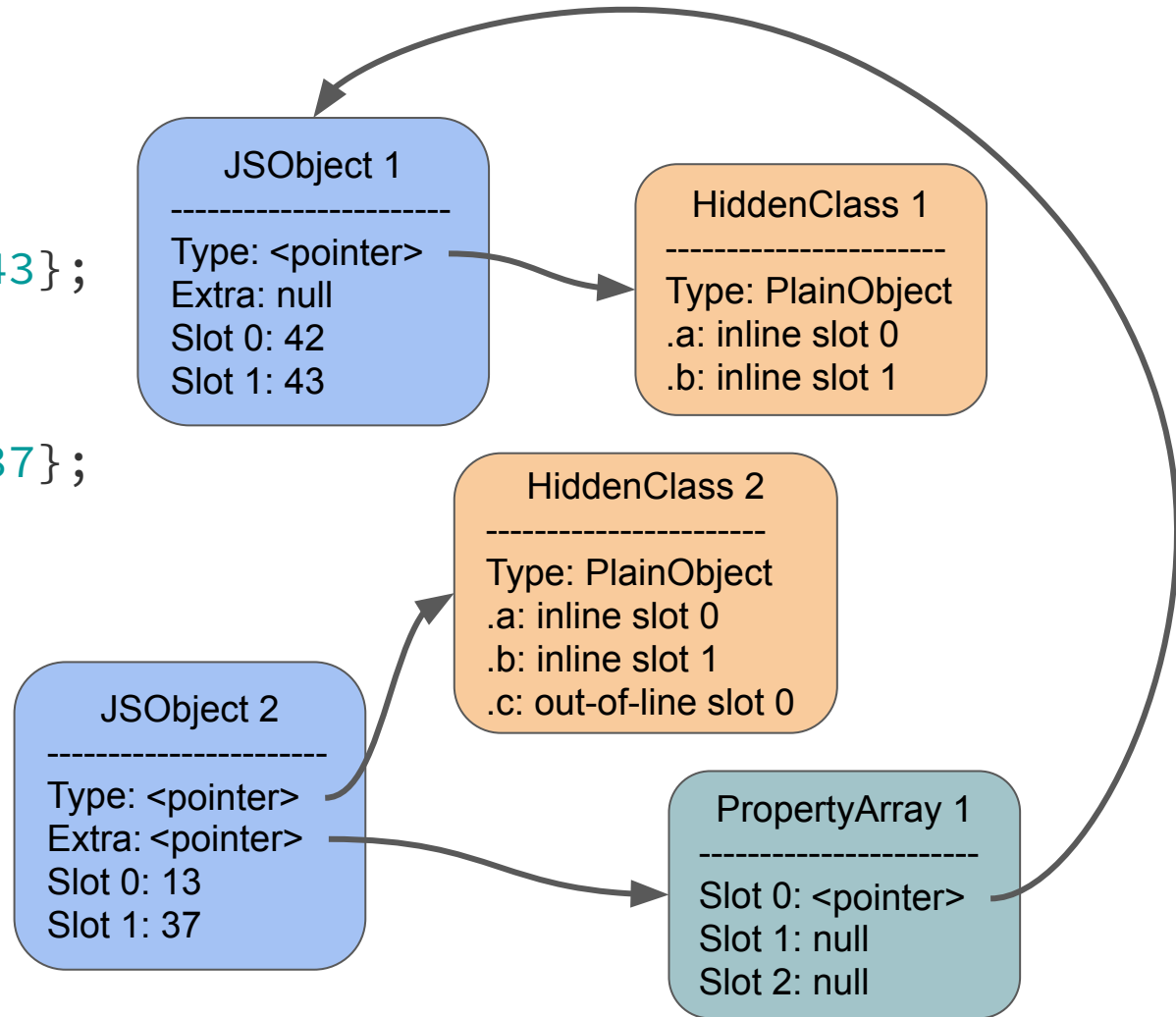
Basic JavaScript

```
let o1 = {a: 42, b: 43};  
console.log(o1.a);  
  
let o2 = {a: 13, b: 37};  
  
o2.c = o1;
```



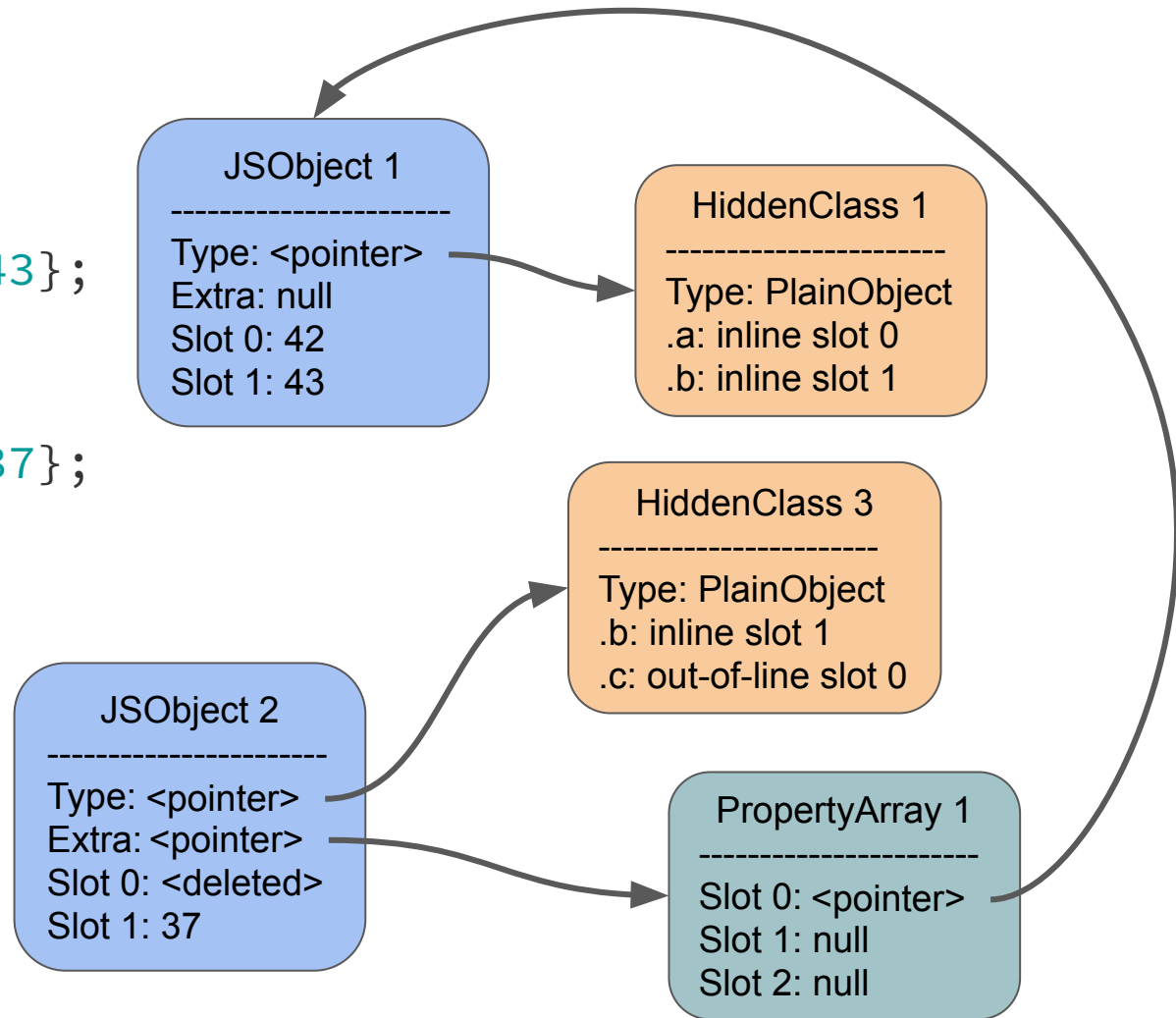
Basic JavaScript

```
let o1 = {a: 42, b: 43};  
console.log(o1.a);  
  
let o2 = {a: 13, b: 37};  
  
o2.c = o1;
```



Basic JavaScript

```
let o1 = {a: 42, b: 43};  
console.log(o1.a);  
  
let o2 = {a: 13, b: 37};  
  
o2.c = o1;  
  
delete o2.a;
```



Bytecode
Compiler

Interpreter

Runtime

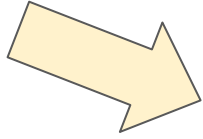
(objects, globals, constructors,
functions, methods, ...)

JIT
Compiler(s)

Wasm
Compiler(s)

Garbage
Collector
(GC)

```
function main() {  
  console.log("Hello World!");  
}  
main();
```



Bytecode
Compiler

Interpreter

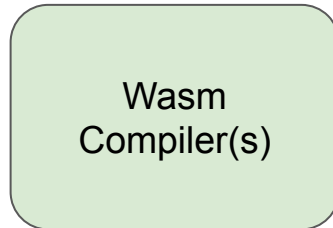
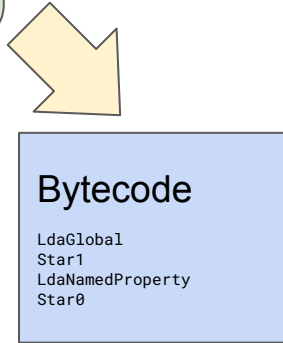
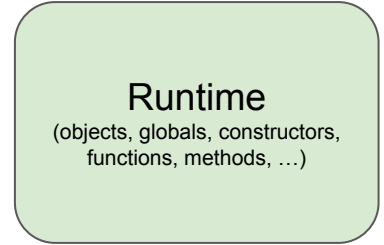
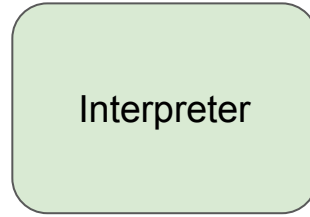
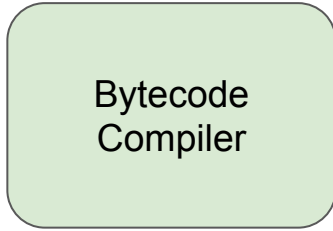
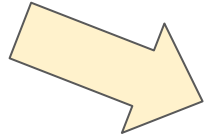
Runtime
(objects, globals, constructors,
functions, methods, ...)

JIT
Compiler(s)

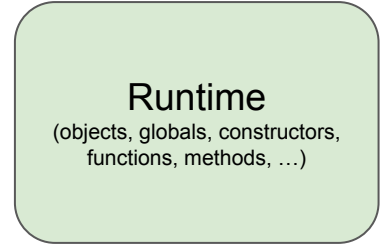
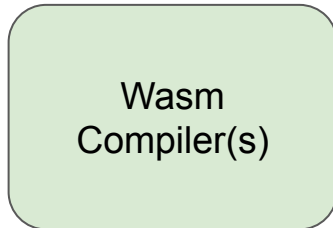
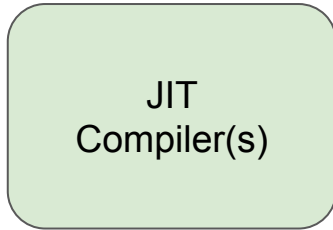
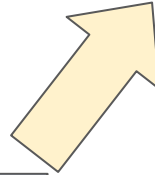
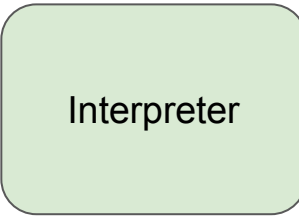
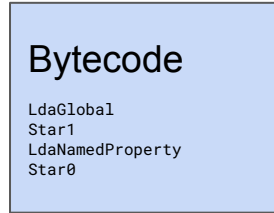
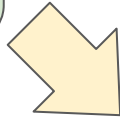
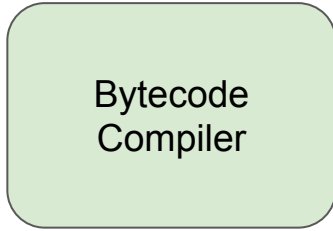
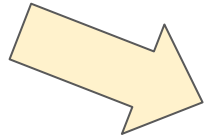
Wasm
Compiler(s)

Garbage
Collector
(GC)

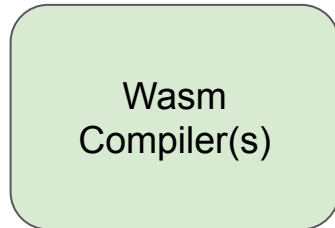
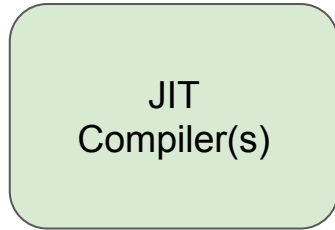
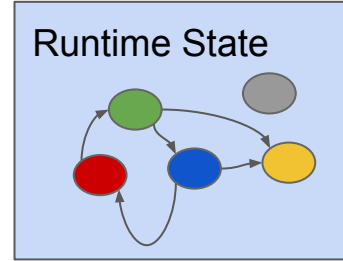
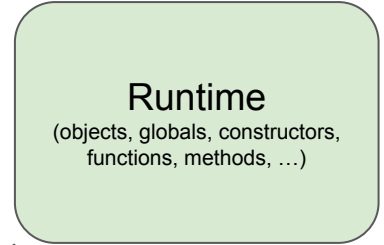
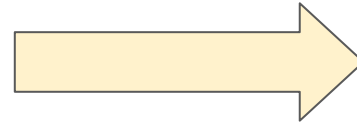
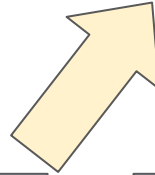
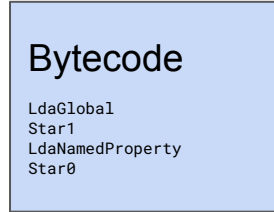
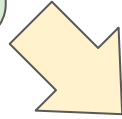
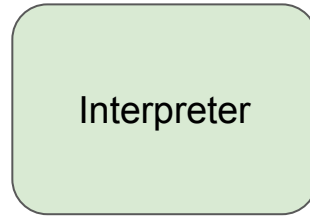
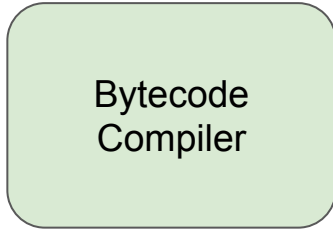
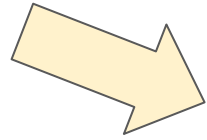
```
function main() {  
  console.log("Hello World!");  
}  
main();
```

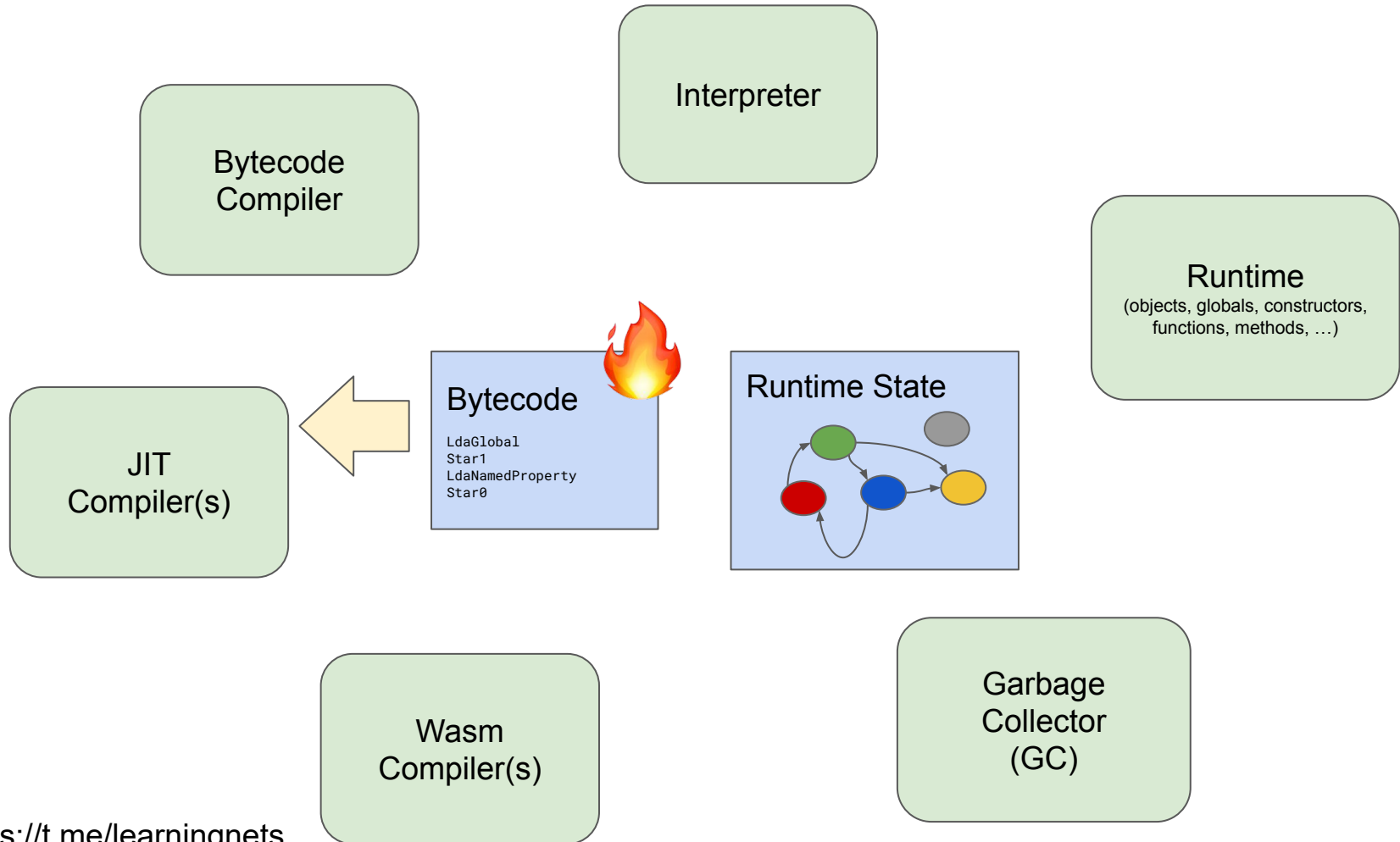


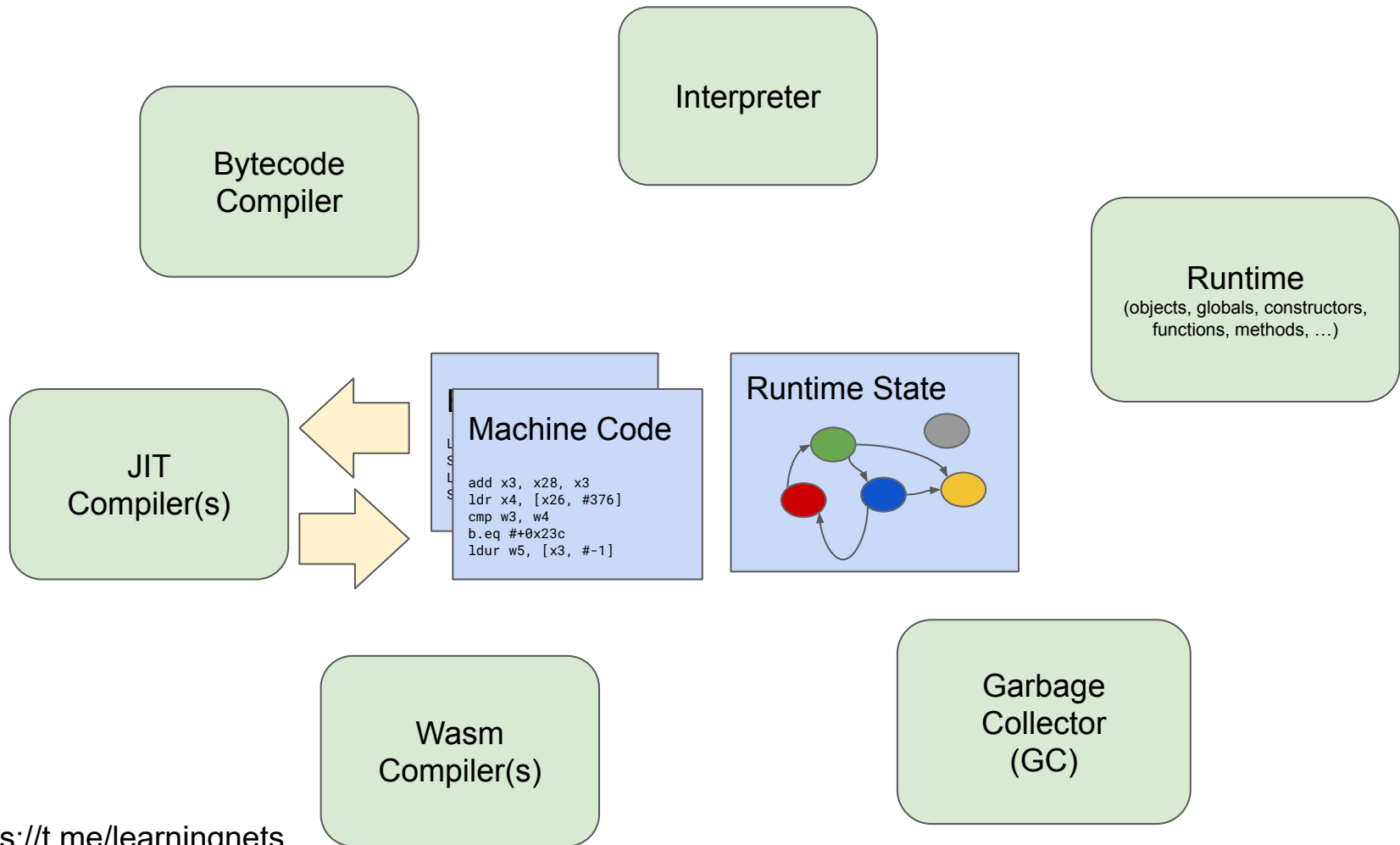
```
function main() {  
  console.log("Hello World!");  
}  
main();
```

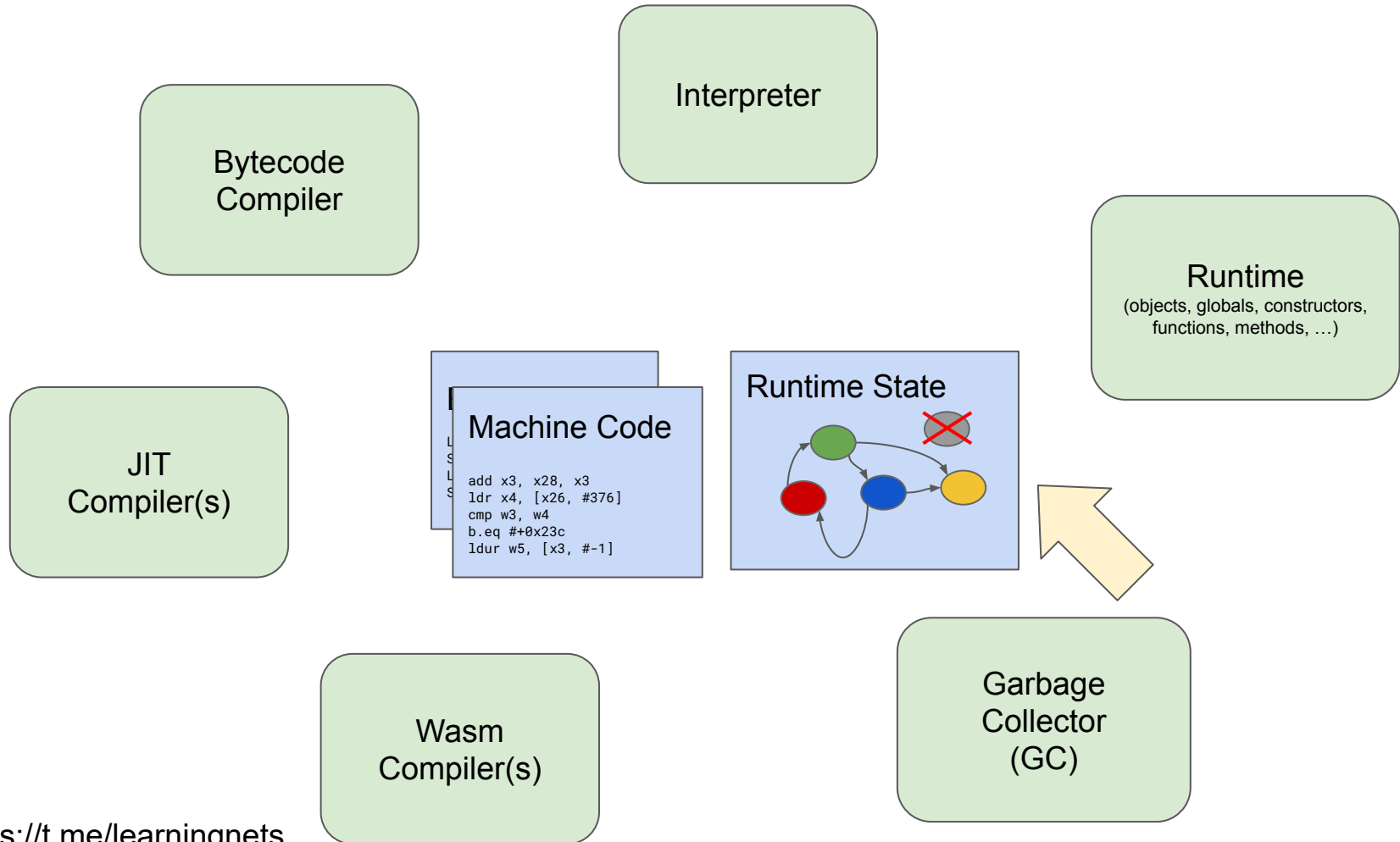


```
function main() {  
  console.log("Hello World!");  
}  
main();
```









Bytecode
Compiler

Interpreter

Runtime
(objects, globals, constructors,
functions, methods, ...)

JIT
Compiler(s)

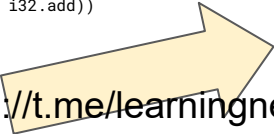
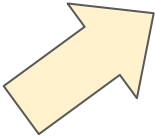
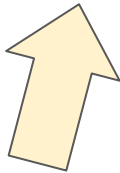
Bytecode
LdaGlobal
Star1
LdaNamedProperty
Star0

Machine Code
add x3, x28, x3
ldr x4, [x26, #376]
cmp w3, w4
b.eq #+0x23c
ldur w5, [x3, #-1]

Wasm
Compiler(s)

Garbage
Collector
(GC)

```
(func  
  (param $lhs i32)  
  (param $rhs i32)  
  (result i32)  
    local.get $lhs  
    local.get $rhs  
    i32.add))
```



JIT Compilation

A (Hypothetical) JIT Optimization Example

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
    if (p.x < 0 || p.x >= W ||
        p.y < 0 || p.y >= H) {
        throw "invalid point";
    }
    bmp[p.x * W + p.y] = v;
}
```

Example: “Training” the JIT

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
    if (p.x < 0 || p.x >= W ||
        p.y < 0 || p.y >= H) {
        throw "invalid point";
    }
    bmp[p.x * W + p.y] = v;
}

// "Train" the JIT
for (let i = 0; i < 10000; i++) {
    set({x: 1, y: 2}, 3);
}
```

Example: Bytecode Parsing

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
  if (p.x < 0 || p.x >= W ||
      p.y < 0 || p.y >= H) {
    throw "invalid point";
  }
  bmp[p.x * W + p.y] = v;
}
```

```
x1 = LoadProperty p, 'x'
GotoIf .throwException, x1 < 0
x2 = LoadProperty p, 'x'
GotoIf .throwException, x2 >= 64
```

Example: Speculation + Lowering

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
  if (p.x < 0 || p.x >= W ||
      p.y < 0 || p.y >= H) {
    throw "invalid point";
  }
  bmp[p.x * W + p.y] = v;
}
```

```
CheckType p, ObjType1
x1 = LoadField p, +8
GotoIf .throwException, x1 < 0
CheckType p, ObjType1
x2 = LoadField p, +8
GotoIf .throwException, x2 >= 64
```

Example: Speculation + Lowering

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
  if (p.x < 0 || p.x >= W ||
      p.y < 0 || p.y >= H) {
    throw "invalid point";
  }
  bmp[p.x * W + p.y] = v;
}
```

```
CheckType p, ObjType1
x1 = LoadField p, +8
GotoIf .throwException, x1 < 0
CheckType p, ObjType1
x2 = LoadField p, +8
GotoIf .throwException, x2 >= 64
```

Example: Redundancy Elimination

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
  if (p.x < 0 || p.x >= W ||
      p.y < 0 || p.y >= H) {
    throw "invalid point";
  }
  bmp[p.x * W + p.y] = v;
}
```

```
CheckType p, ObjType1
x1 = LoadField p, +8
GotoIf .throwException, x1 < 0
CheckType p, ObjType1
x2 = LoadField p, +8
GotoIf .throwException, x1 >= 64
```

Example: Bytecode Parsing

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
    if (p.x < 0 || p.x >= W ||
        p.y < 0 || p.y >= H) {
        throw "invalid point";
    }
    bmp[p.x * W + p.y] = v;
}
```

```
W = LoadGlobal 'W'
i1 = Mul x, W
i2 = Add i1, y
bmp = LoadGlobal 'bmp'
StoreElement bmp, i2, v
```

Example: Constant Folding + Lowering

```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
    if (p.x < 0 || p.x >= W ||
        p.y < 0 || p.y >= H) {
        throw "invalid point";
    }
    bmp[p.x * W + p.y] = v;
}
```

```
i1 = IntegerMul x, 64
i2 = IntegerAdd i1, y
CheckBounds i2, 4096
CheckType v, Uint8
StoreUint8Array bmp, i2, v
```

Example: Range Analysis + Bounds Check Elimination

```
const W = 64, H = 64; // x = Range [0, 64)
const bmp = new Uint8Array(W * H); // y = Range [0, 64)
function set(p, v) {
    if (p.x < 0 || p.x >= W ||
        p.y < 0 || p.y >= H) {
        throw "invalid point";
    }
    bmp[p.x * W + p.y] = v;
}
// i1 = IntegerMul x, 64
// i1 = Range [0, 4033)
i2 = IntegerAdd i1, y
// i2 = Range [0, 4096)
CheckBounds i2, 4096
...
```

Example: Final JIT IR Code

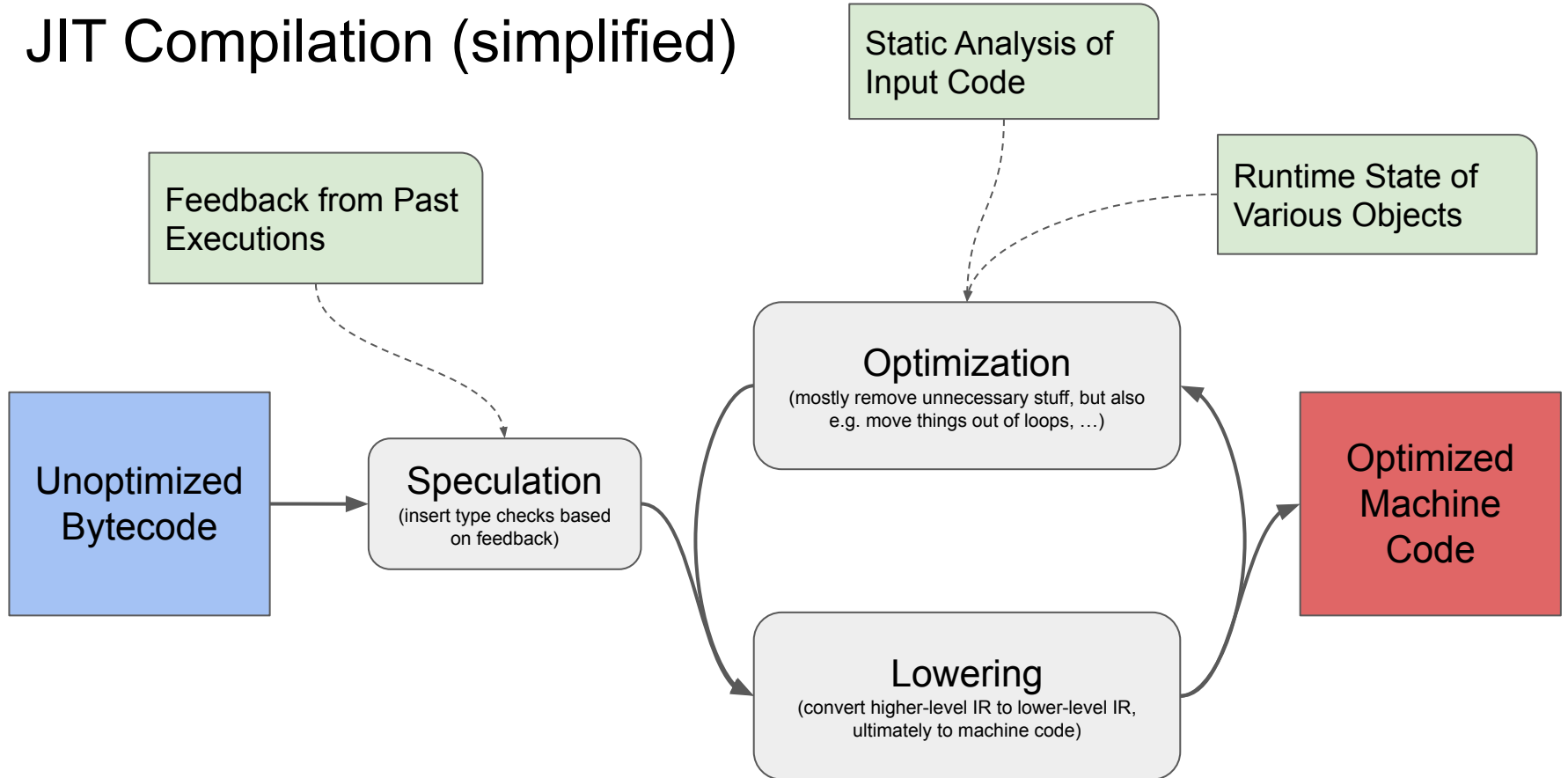
```
const W = 64, H = 64;
const bmp = new Uint8Array(W * H);
function set(p, v) {
  if (p.x < 0 || p.x >= W ||
      p.y < 0 || p.y >= H) {
    throw "invalid point";
  }
  bmp[p.x * W + p.y] = v;
}
```

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```
CheckType p, ObjType1
x = LoadField p, +8
y = LoadField p, +16
GotoIf .throwException x < 0 || ...
```

```
i1 = IntegerMul x, 64
i2 = IntegerAdd i1, y
CheckType v, Uint8
StoreUint8Array bmp, i2, v
```

JIT Compilation (simplified)



A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
    let i = a.findIndex(cond);  
    a[i] = v;  
}
```

```
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => e == 3, 42);  
  
// a == [0, 1, 2, 42, 4, 5];
```

Description

The `findIndex()` method executes the `callbackFn` function once for every index in the array until it finds the one where `callbackFn` returns a [truthy](#) value.

A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
    let i = a.findIndex(cond);  
    a[i] = v;  
}  
  
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => e == 3, 42);  
  
// a == [0, 1, 2, 42, 4, 5];
```

CheckType a, ArrType1
i = Call Runtime_FindIndex(a, cond)
CheckBounds a, i
StoreArray a, i, v

A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
  let i = a.findIndex(cond);  
  a[i] = v;  
}
```

```
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => e == 3, 42);  
// a == [0, 1, 2, 42, 4, 5];
```

```
CheckType a, ArrType1  
i = Call Runtime_FindIndex(a, cond)  
// i = Range [0, a.length - 1)  
CheckBounds a, i  
StoreArray a, i, v
```

A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
    let i = a.findIndex(cond);  
    a[i] = v;  
}  
  
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => false, 42);
```

CheckType a, ArrType1
i = Call Runtime_FindIndex(a, cond)
// i = Range [0, a.length - 1)
~~CheckBounds a, i~~
StoreArray a, i, v

Return value

The index of the first element in the array that passes the test. Otherwise, `-1`.

A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
  let i = a.findIndex(cond);  
  a[i] = v;  
}  
  
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => false, 42);
```

```
CheckType a, ArrType1  
i = Call Runtime_FindIndex(a, cond)  
// i = Range [0, a.length - 1)  
CheckBounds a, i  
StoreArray a, i, v
```

Return value

The index of the first element in the array that passes the test. Otherwise, `-1`.

A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
    let i = a.findIndex(cond);  
    a[i] = v;  
}  
  
CheckType a, ArrType1  
i = Call Runtime_FindIndex(a, cond)  
// i = Range [-1, a.length - 1)  
Check i >= 0  
StoreArray a, i, v
```

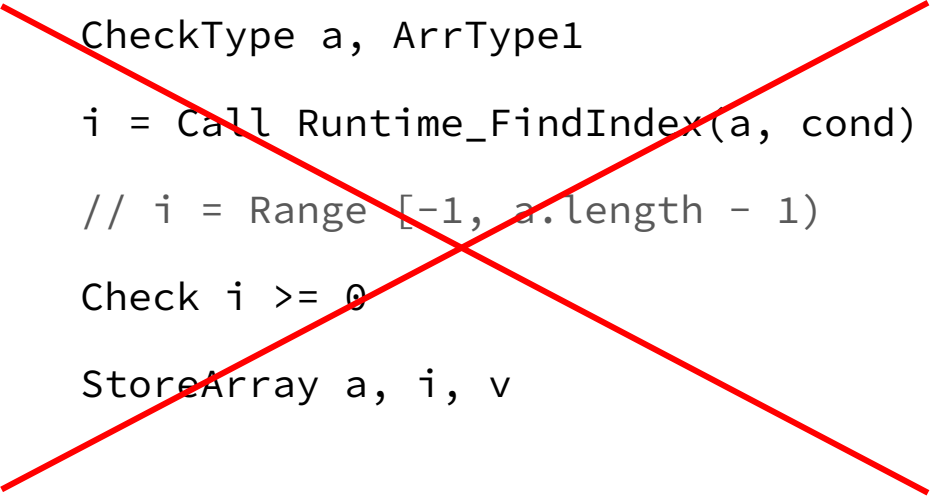
A (Hypothetical) JIT Bug Example

```
function replace(a, cond, v) {  
    let i = a.findIndex(cond);  
    a[i] = v;  
}  
  
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => {  
    a.length = 0; return true;  
}, 42);
```

CheckType a, ArrType1
i = Call Runtime_FindIndex(a, cond)
// i = Range [-1, a.length - 1)
Check i >= 0
StoreArray a, i, v

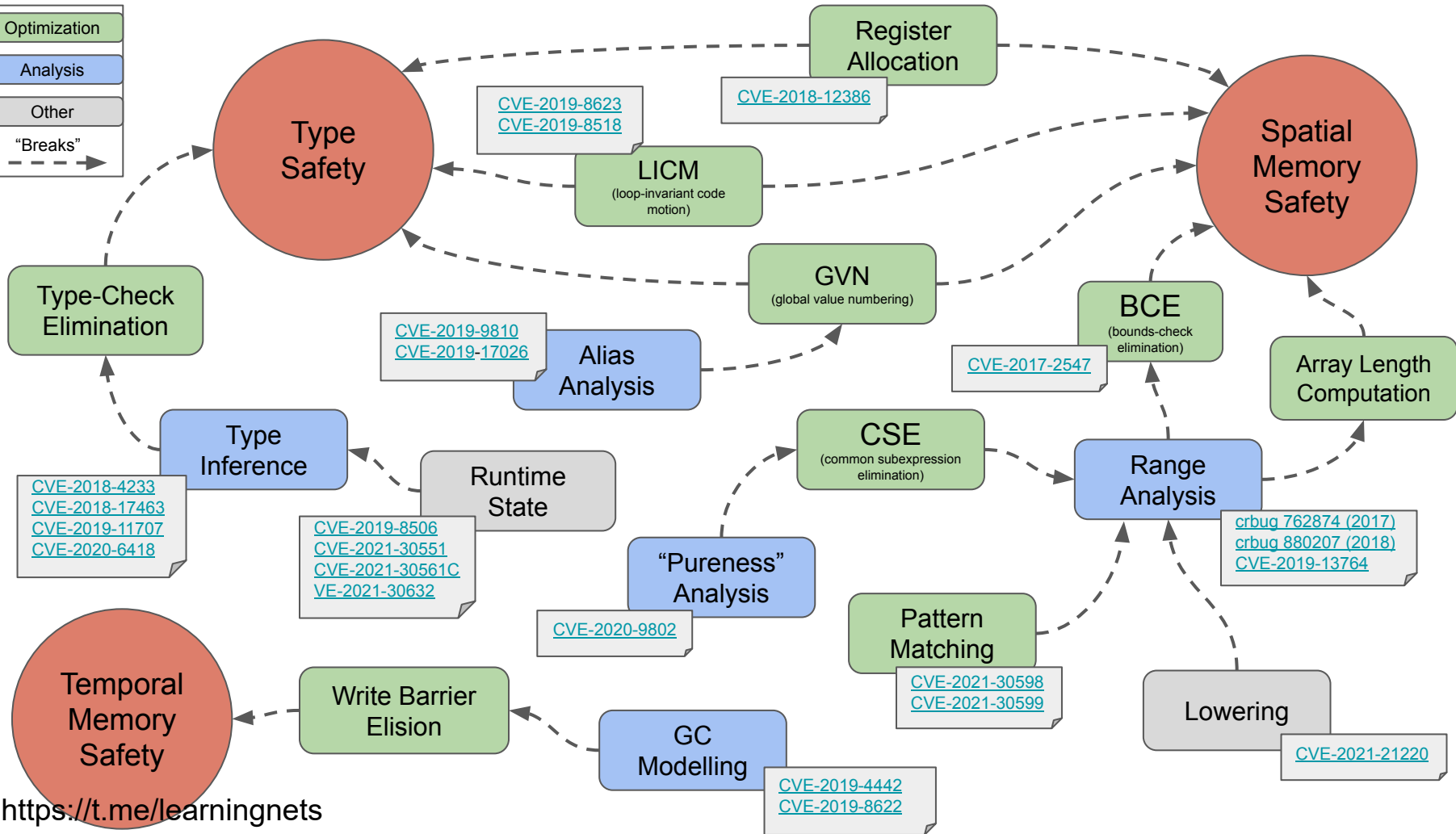
A (Hypothetical) JIT Bug Example

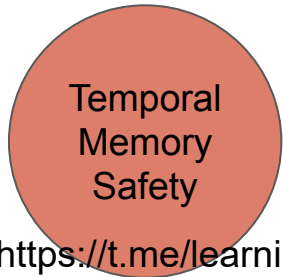
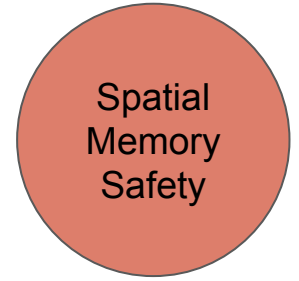
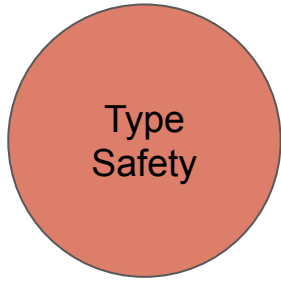
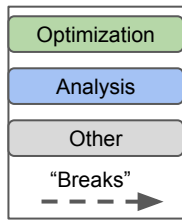
```
function replace(a, cond, v) {  
  let i = a.findIndex(cond);  
  a[i] = v;  
}  
  
let a = [0, 1, 2, 3, 4, 5];  
replace(a, (e) => {  
  a.length = 0; return true;  
}, 42);
```

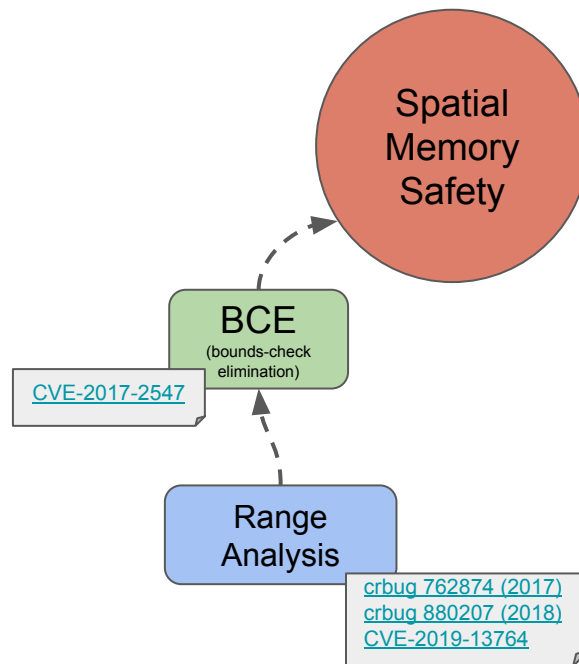
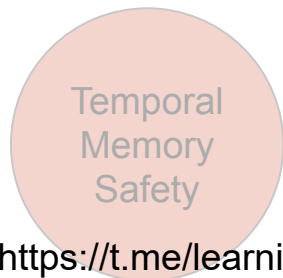
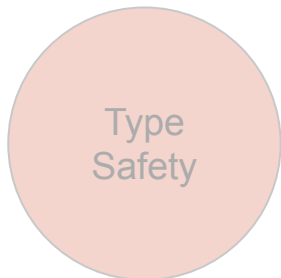
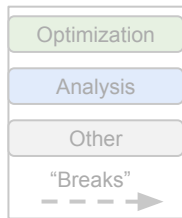


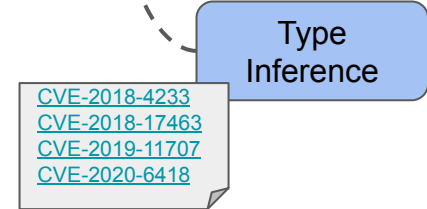
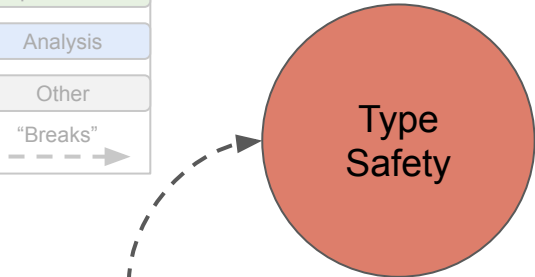
```
CheckType a, ArrType1  
i = Call Runtime_FindIndex(a, cond)  
// i = Range [-1, a.length - 1)  
Check i >= 0  
StoreArray a, i, v
```

Optimization
Analysis
Other
"Breaks" →



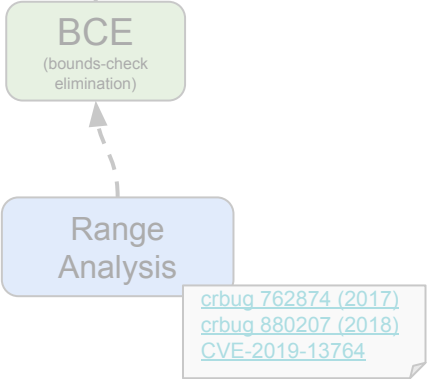
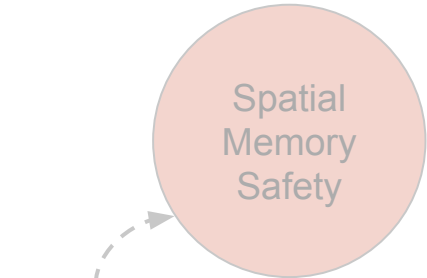
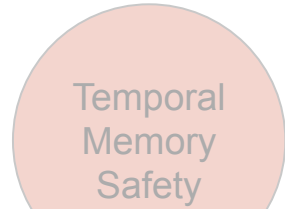




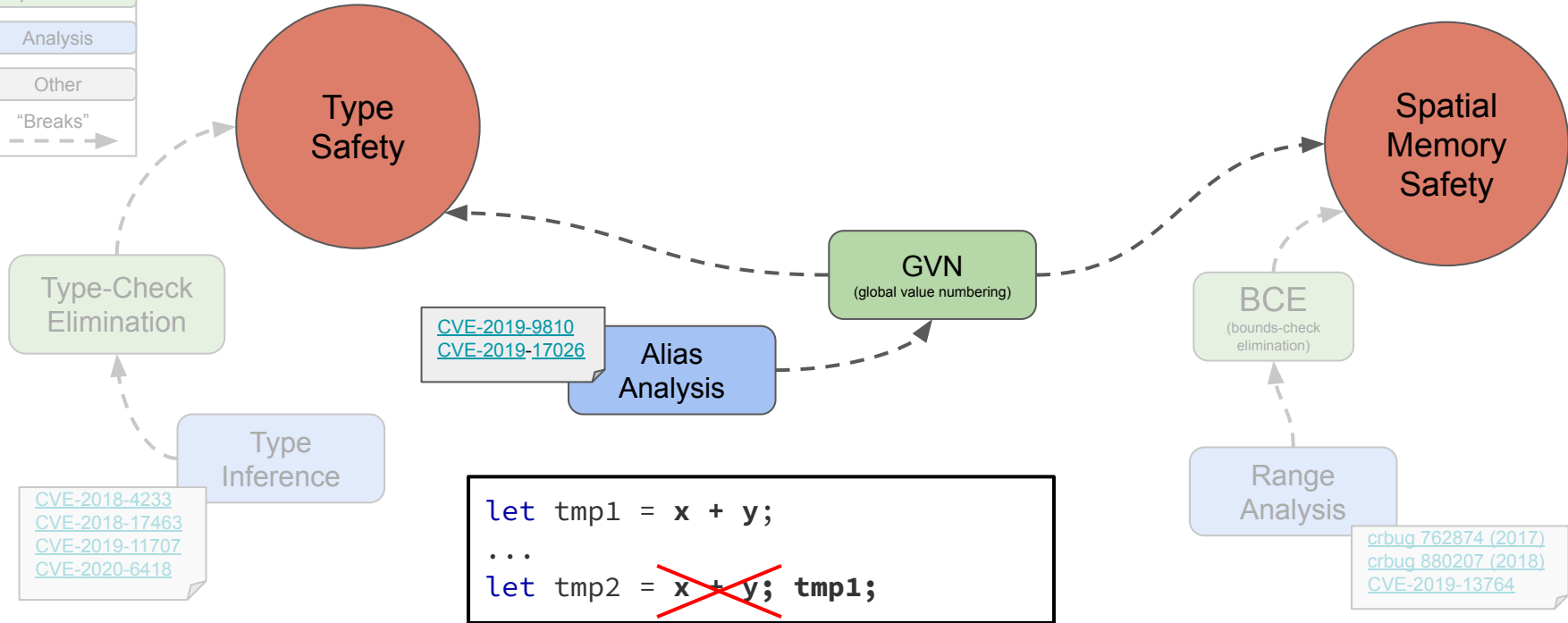


```

CheckType o, ObjType1
...
CheckType o, ObjType1
    
```



Optimization
Analysis
Other
“Breaks”



[CVE-2019-9810](#)
[CVE-2019-17026](#)

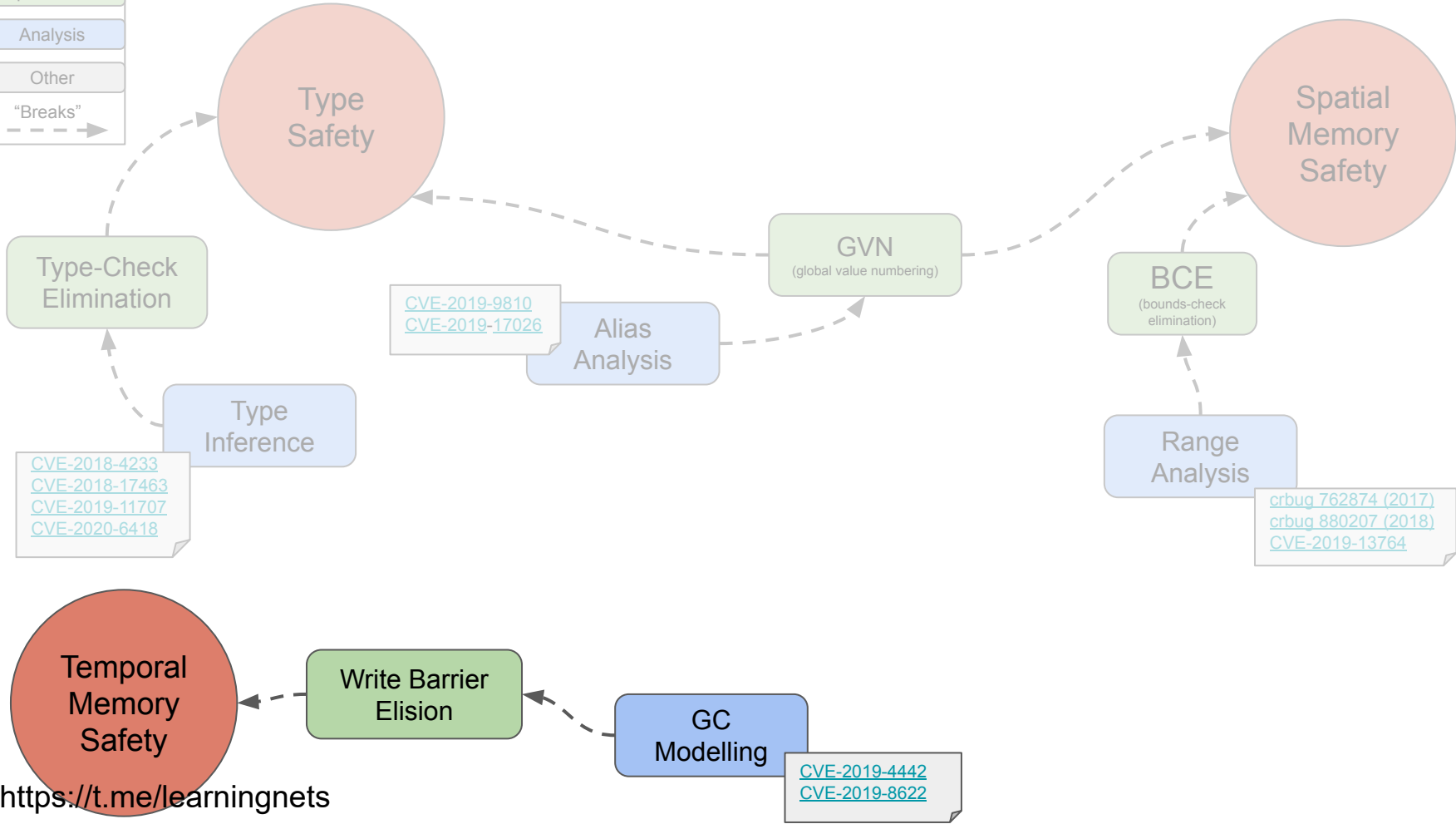
[CVE-2018-4233](#)
[CVE-2018-17463](#)
[CVE-2019-11707](#)
[CVE-2020-6418](#)

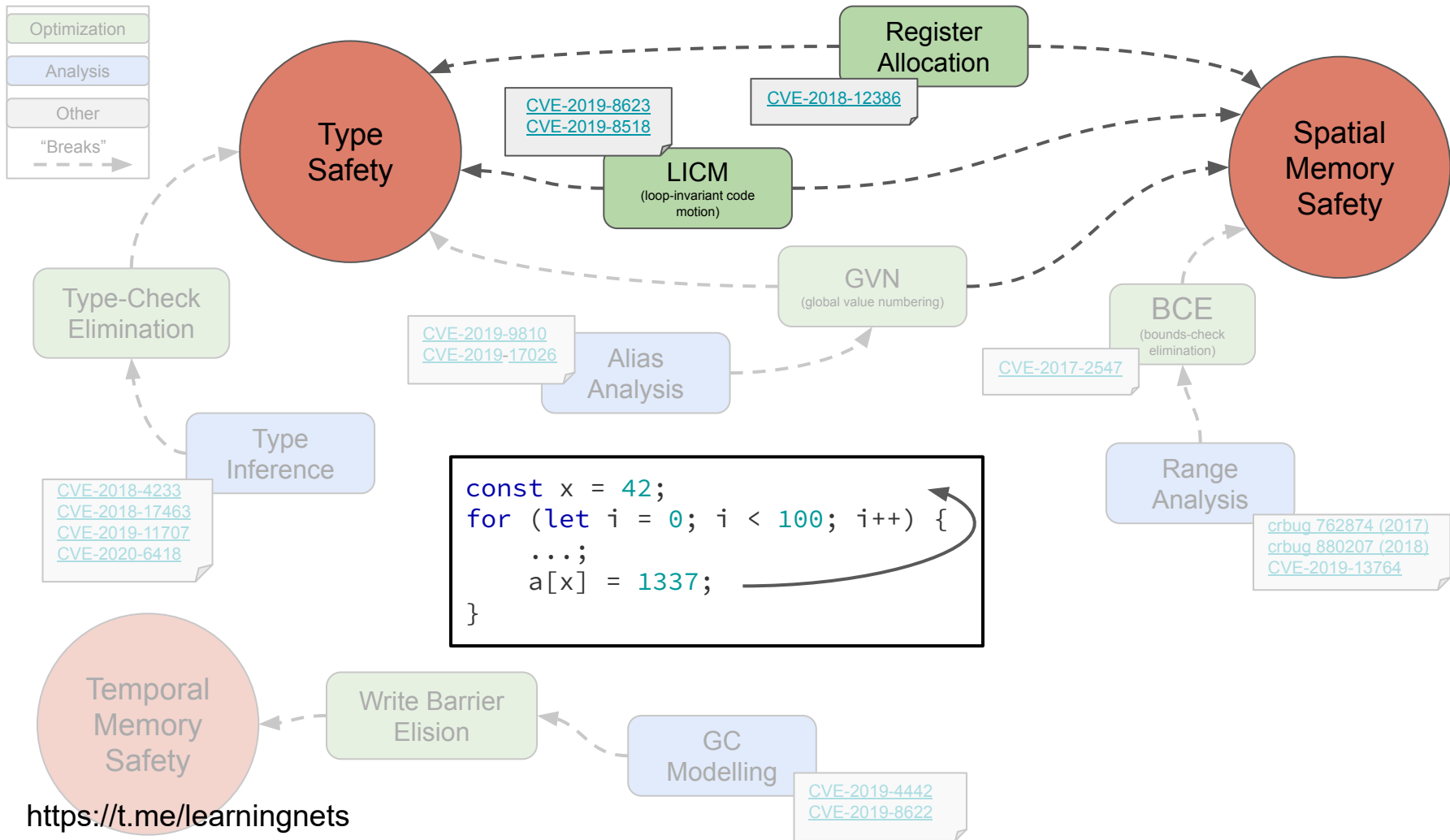
```
let tmp1 = x + y;
...
let tmp2 = x + y; tmp1;
```

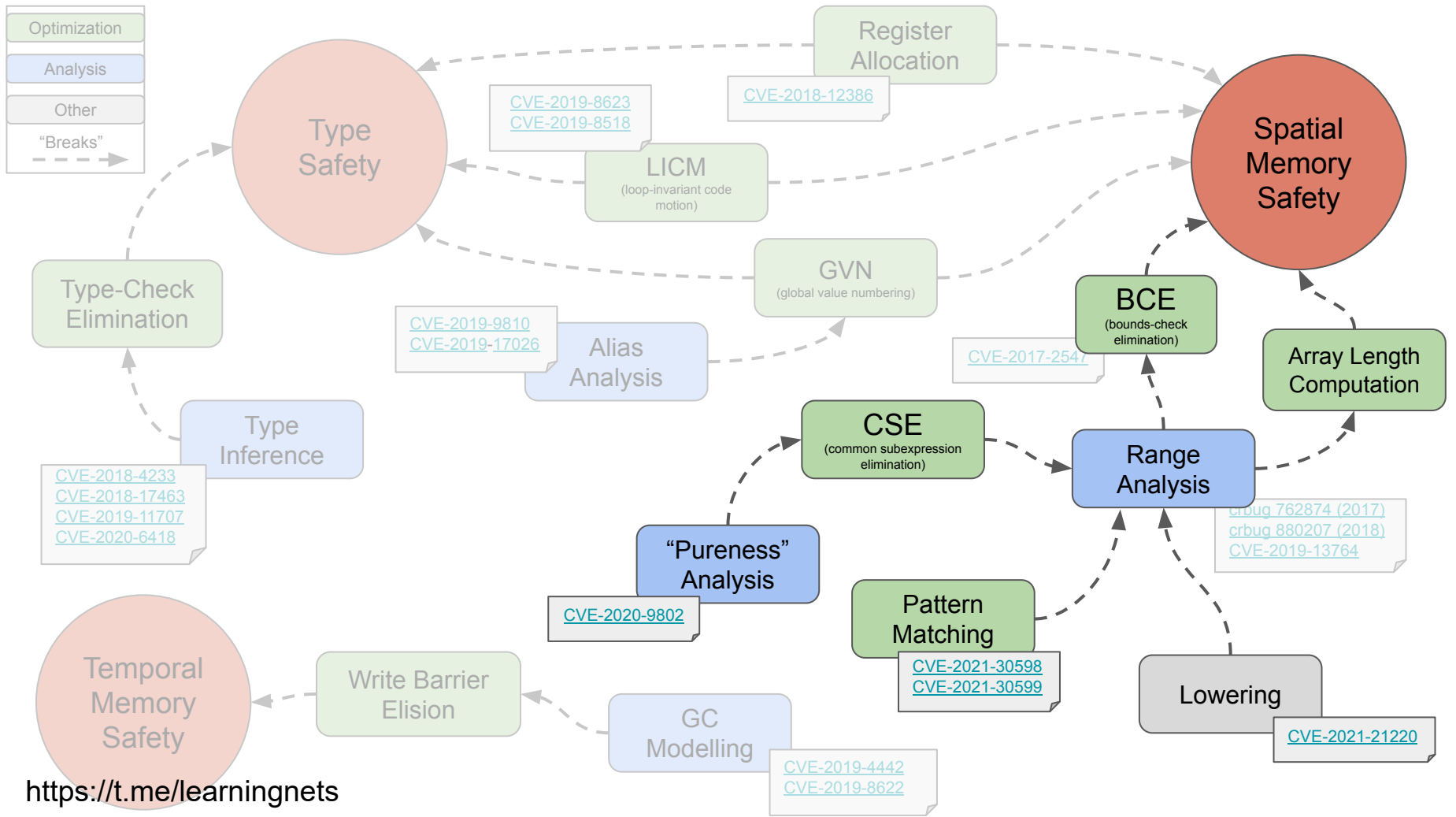
[crbug 762874 \(2017\)](#)
[crbug 880207 \(2018\)](#)
[CVE-2019-13764](#)

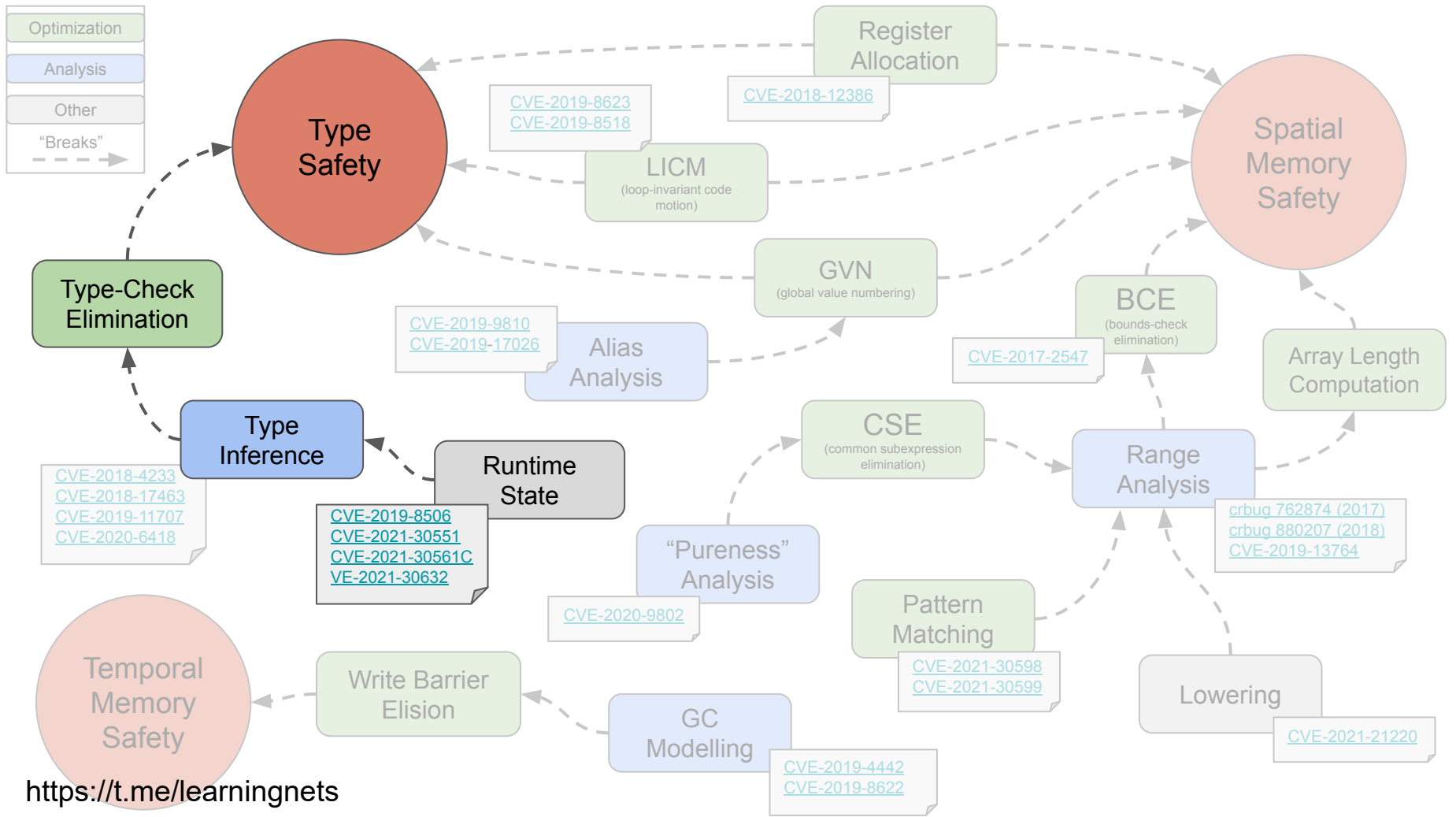
Temporal
Memory
Safety

Optimization
Analysis
Other
“Breaks” ----->

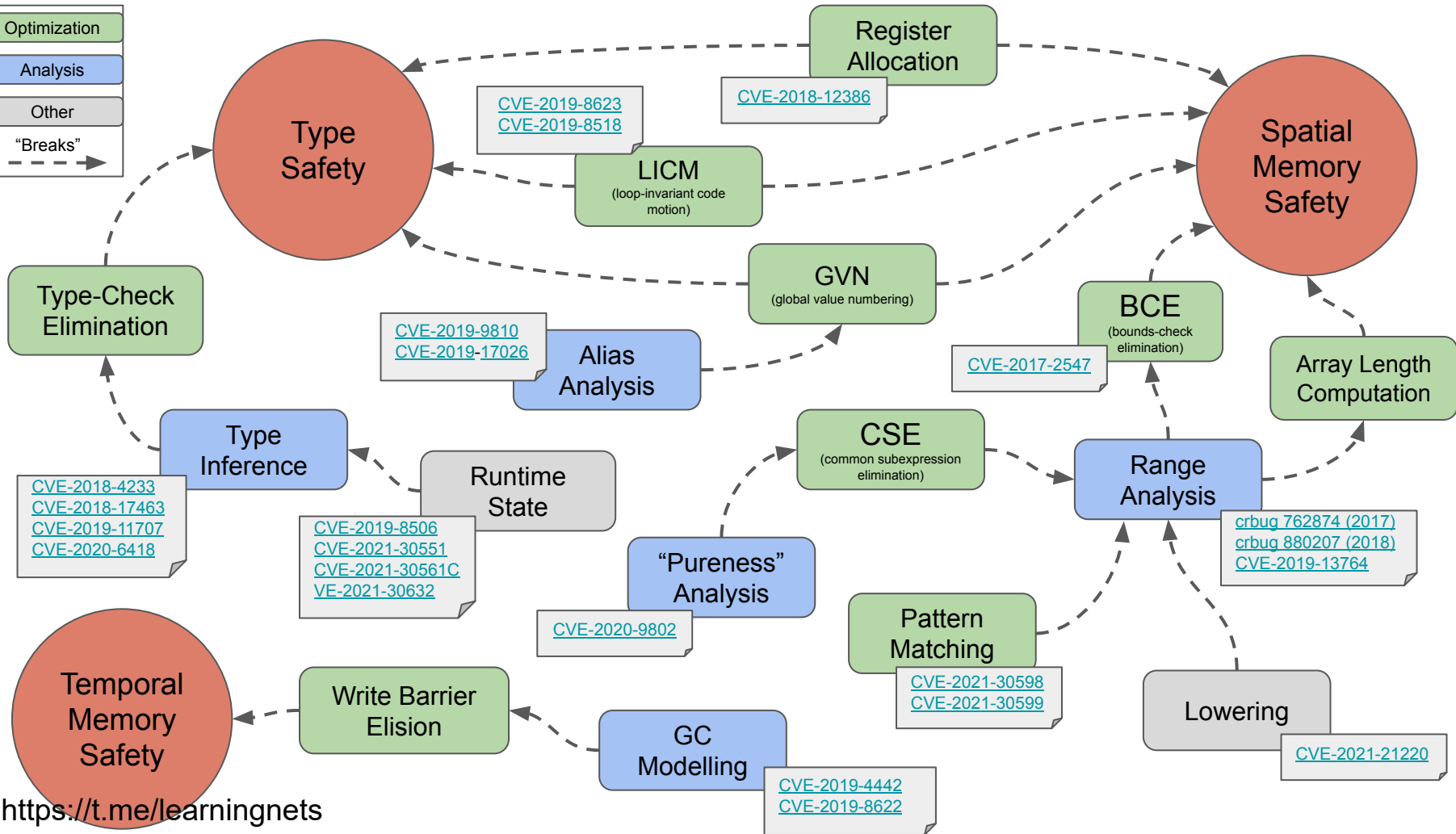


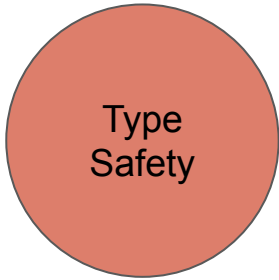






Optimization
Analysis
Other
"Breaks" →

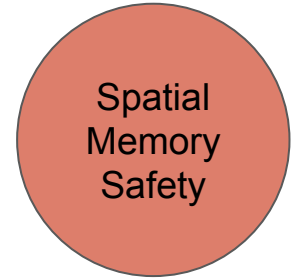




Type
Safety

Exploitation

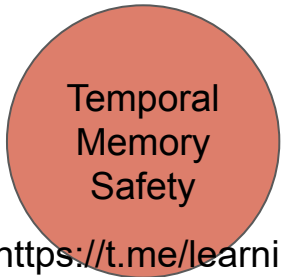
- Choose (arbitrary) victim type
- Choose (arbitrary) target type
- Choose (arbitrary) operation
- Trigger bug to confuse objects



Spatial
Memory
Safety

Exploitation

- Choose (arbitrary) victim array
- Choose (arbitrary) OOB index
- Choose read or write access
- Trigger bug to corrupt memory



Temporal
Memory
Safety

Exploitation

- Choose (arbitrary) victim type
- Choose (arbitrary) replacement type
- Trigger bug and GC to cause UaF

JS Outside JIT

Bytecode
Compiler

[CVE-2022-0102](#)

Interpreter

[CVE-2021-30517](#)
[CVE-2021-38001](#)

Runtime

(objects, globals, constructors,
functions, methods, ...)

[CVE-2021-1789](#)
[CVE-2021-21225](#)
[CVE-2021-38003](#)

JIT
Compiler(s)

See prev. slides :)

- Plenty of complexity elsewhere
- Few bug patterns, many “1-off” bugs

Wasm
Compiler(s)

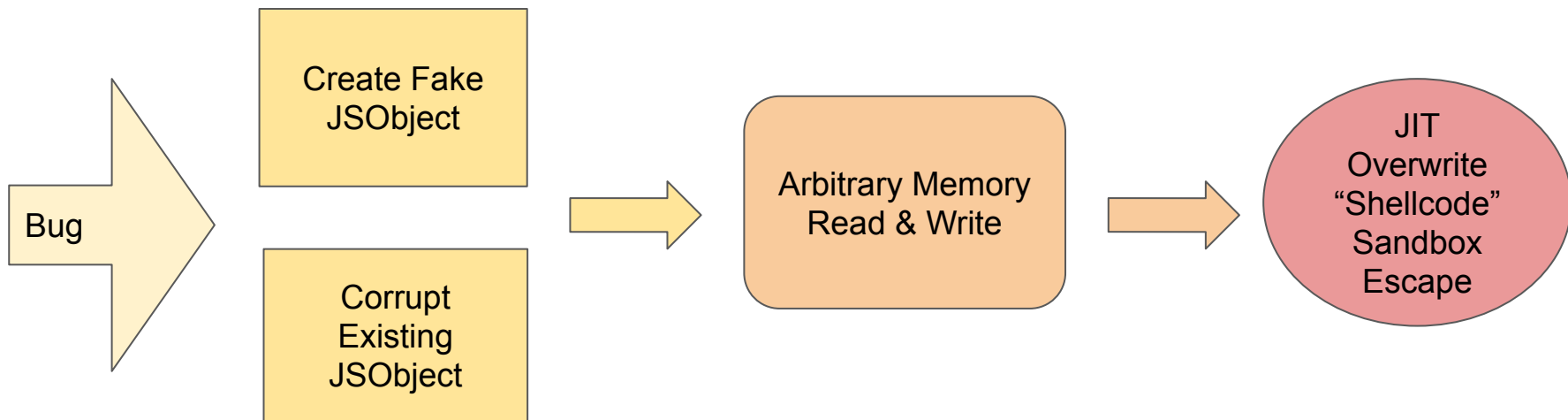
[CVE-2021-30734](#)

Garbage
Collector
(GC)

[CVE-2021-37975](#)

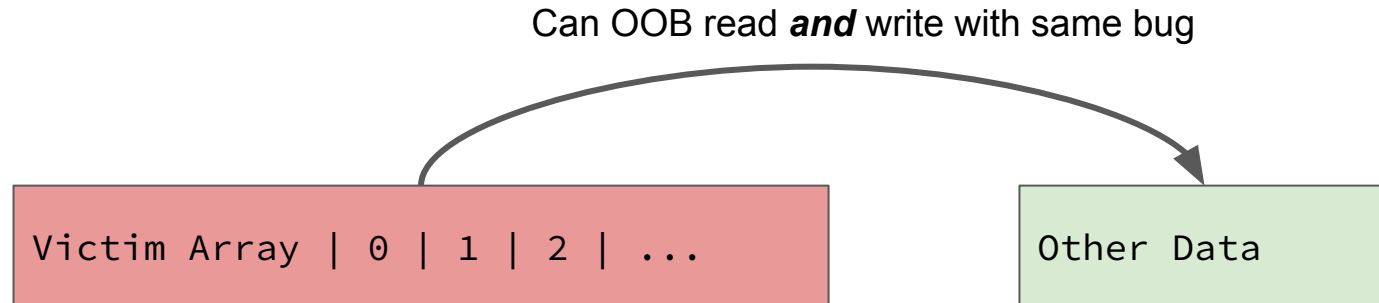
Exploitation & Mitigations

Exploit Flow Circa 2016



What About Classical Mitigations?

- ASLR: Usually easy to construct a leak via type confusion or OOB (second bug not required)
- DEP/NX: JIT provides easy ways to map shellcode
- Stack Cookies: Most JS bugs are heap based



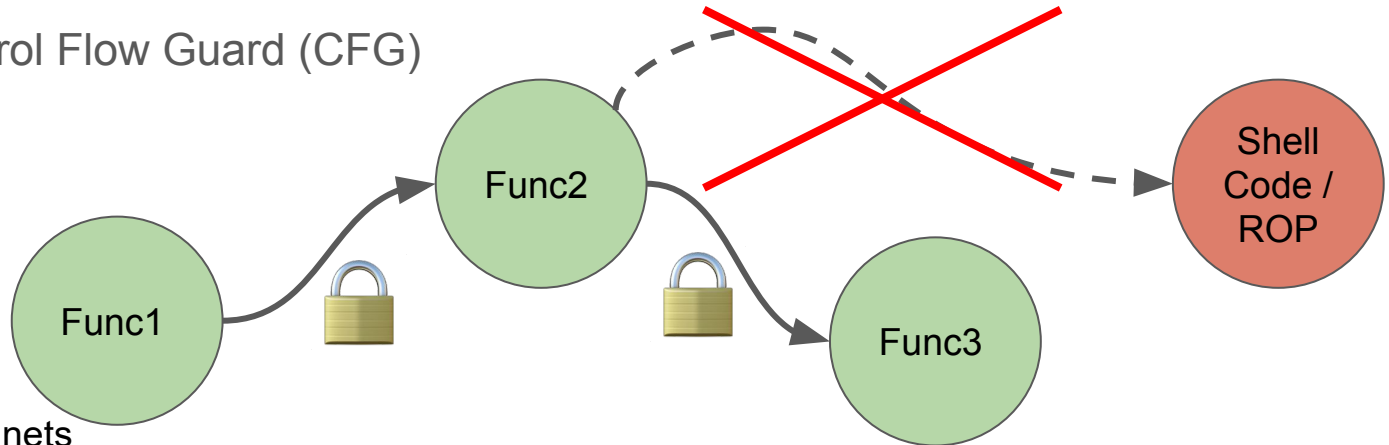
"Modern" Mitigations

When most people think of modern mitigations they think of Control Flow Integrity (CFI)

Armv8.3+: Pointer Authentication (PAC) and Branch Target Identification (BTI)

Intel: Shadow Stack and Control-flow Enforcement Technology (CET)

Windows: Control Flow Guard (CFG)



"Modern" Mitigations - Not Quite There Yet

When most people think of modern mitigations they think of Control Flow Integrity (CFI)

~~Armv8.3+: Pointer Authentication (PAC) and Branch Target Identification (BTI)~~

~~Intel: Shadow Stack and Control flow Enforcement Technology (CET)~~

~~Windows: Control Flow Guard (CFG)~~

JSC supports PAC

V8 does not yet have full support for [CET](#), [CFG](#), or [PAC](#)

Pointer Authentication

Newer iOS devices and M1 Macbooks benefit from Armv8's Pointer Authentication

- PAC*: signs the pointer, writes cryptographic signature to upper bits
- AUT*: verifies the pointer

Mostly used to protect code pointers, but may be used for data as well

0x00007fc75ae25b20



0xa9b67fc75ae25b20

0xa9b6414141414141



0x8000414141414141

Bypassing Pointer Authentication

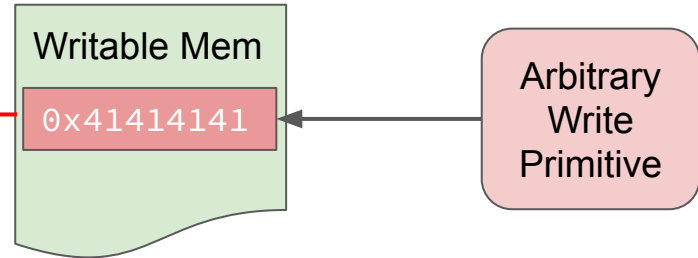
PAC bypasses can be considered similar to bugs; ie patched quickly if disclosed

Example Bypass Methods

- Pointer Forgery: Writable memory which later gets signed [\[ref\]](#)
- Swap or use signed pointers which lack context

```
ADRP    X16, #_pow_ptr_3@PAGE
LDR     X16, [X16,#_pow_ptr_3@PAGEOFF]
PACIZA X16
```

→ 0x66c0000041414141



Additionally, V8 currently supports PAC, but not in JITed code [\[ref\]](#)

"Scripted" Code Execution

If you can't get arbitrary asm code, you may be able to call existing functionality

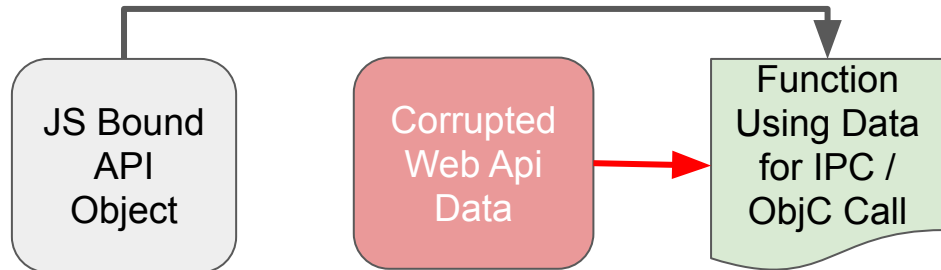
Build control flow with manipulated calls / actions made from JavaScript

Required sandbox escape functionality usually already exists!

Good Example: ObjectiveC Selector Calls [\[ref\]\[ref\]](#)

`corrupted_web_obj.do_action()`

<https://t.me/learningnets>



The Rise Of Data-Only Attacks

On PAC devices and as CFI rolls out, shellcode/rop exec is becoming harder...

However, this is usually not the endgame of a JS exploit

Exploits may attempt to attack cross-process data **integrity** / **confidentiality**

- Corrupt IPC data / messages / state to exploit a sandbox bug
- Read sensitive data stored within the process itself

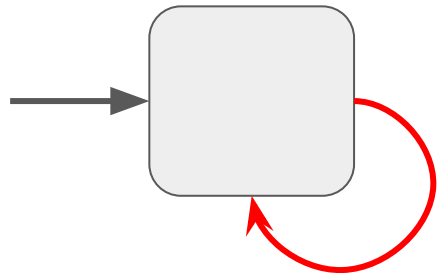
These attacks do not rely on code exec, **only memory read and write**

Exploitation Tricks: Winning Races With Linked Lists


A lot of data attacks become races: Either

- You complete the write in time
- You smash some other data and crash...

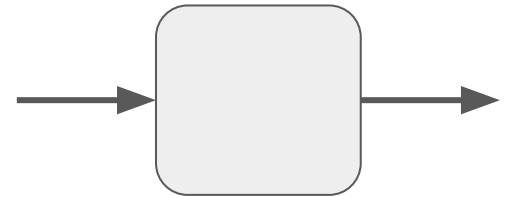
We can abuse linked list structures to stall this race [\[ref\]](#)



Corrupt link list to
form a cycle
<https://t.me/learningsnets>

Wait until
program hits
 loop

Perform
write to
target data

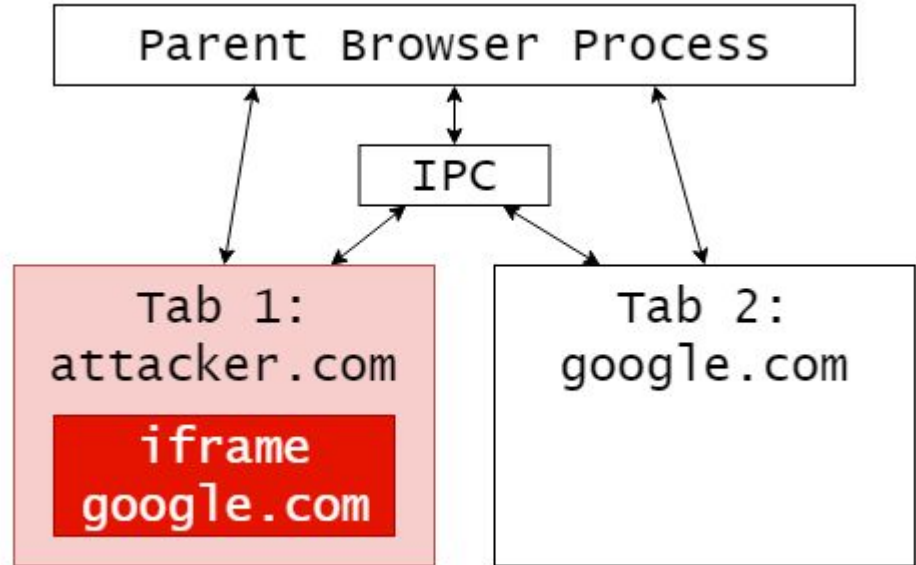


Repair linked list and
allow thread to resume

Attacking Cross-Origin

We have control of all the data in the compromised process

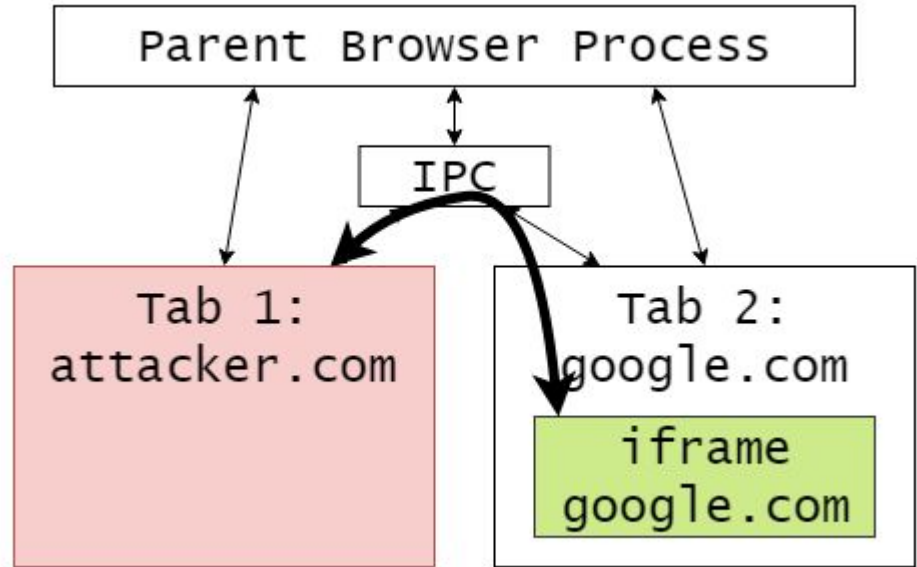
- Force the process to load sensitive data
- Inject JavaScript into other website -> hijack session
- Abuse persistent data features in other websites [\[ref\]](#)



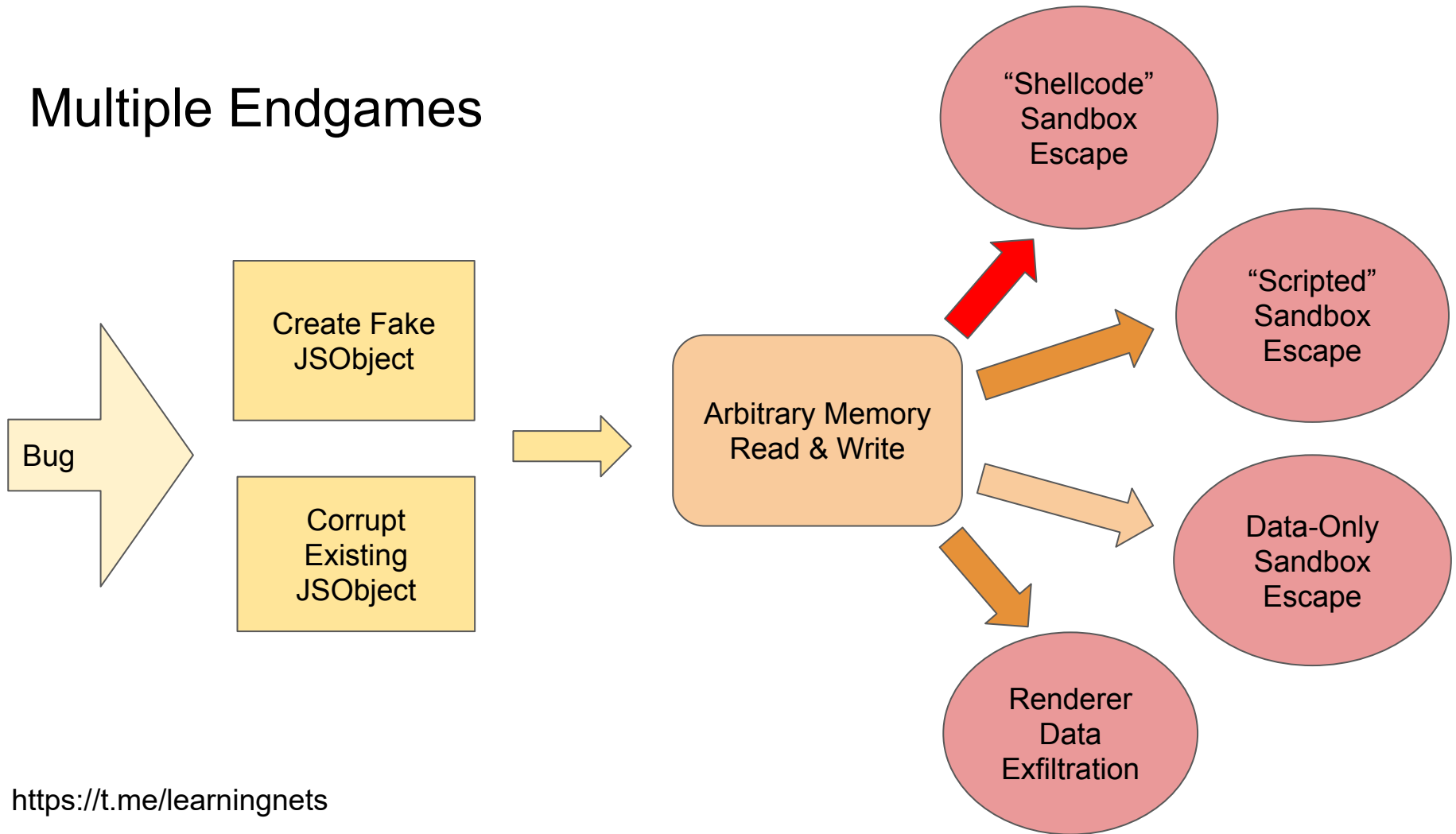
Mitigating Cross Origin Attacks

[Chrome](#) and [Firefox](#) have enabled "Site Isolation"

- Iframes are in separate processes
- Requests and access enforced by the network IPC



Multiple Endgames



Mitigating Arbitrary Read / Write

Arbitrary read/write is a very powerful primitive

Thus, vendors are creating mitigations to make it more difficult

Pointer Caging - Code restricts pointers to specific regions of memory

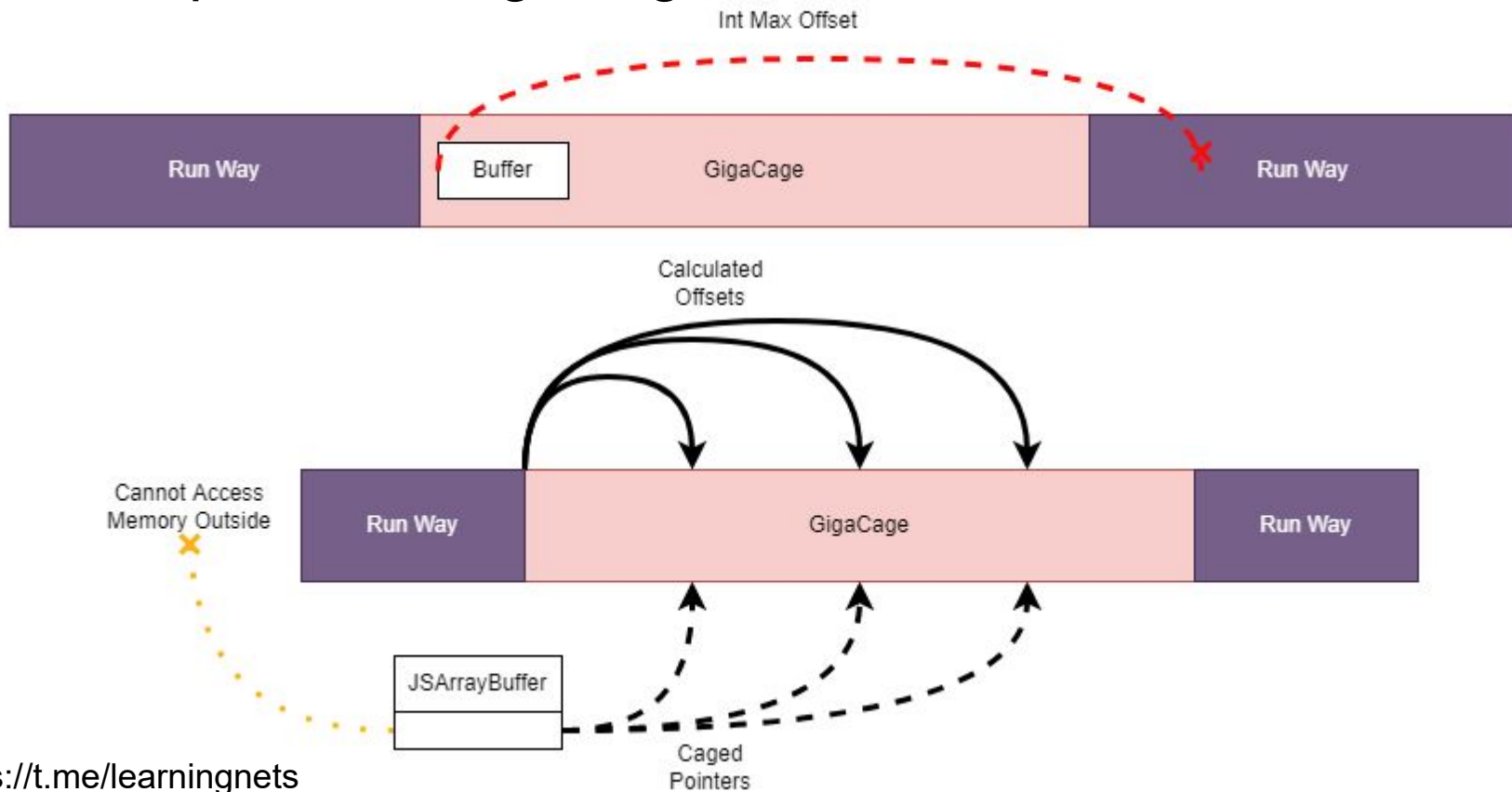
JavaScriptCore's GigaCage

GigaCage prevents pointers being used to corrupt sensitive memory

```
class JSArrayBufferView {  
    using VectorPtr = CagedPtr<Gigacage::Primitive, void, tagCagedPtr>;  
    VectorPtr m_vector;  
}
```

CagedPtr forces all pointer accesses to remain in a specific "GigaCage" region

JavaScriptCore's GigaCage



Is GigaCage Effective?

Required to protect a "vulnerable" pointer:

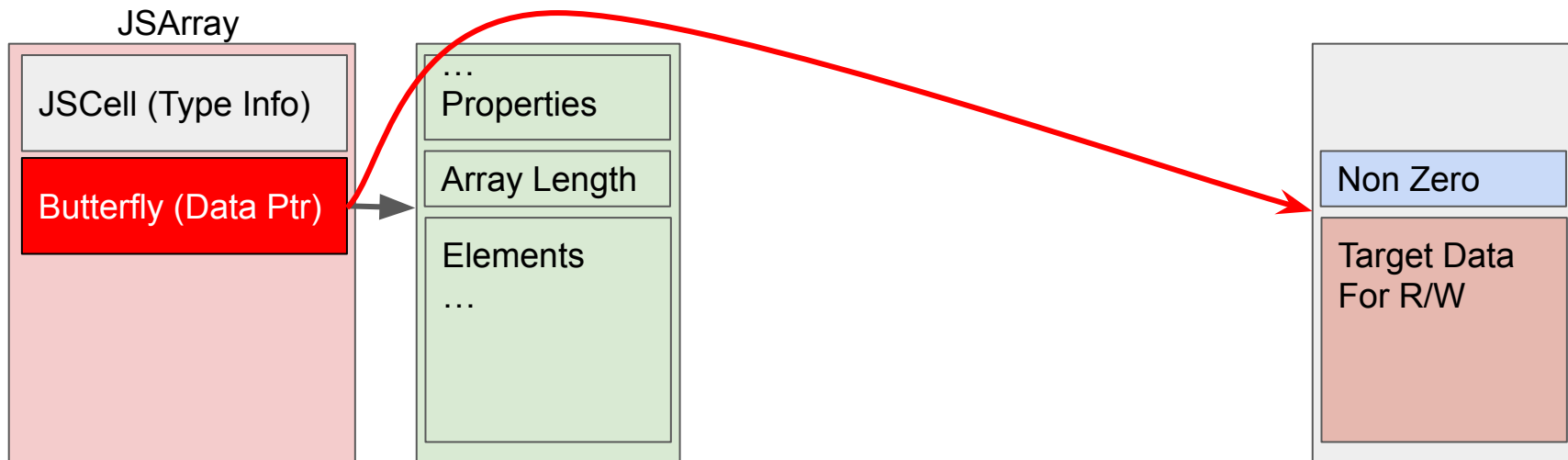
- Explicit caged typing of the pointer
- Correct uncaging implementation when accessing (such as in the JIT)

There are a lot of objects and a lot of pointers

- Attackers just need to find single uncaged pointer they can r/w from
- This is made easier by faking object state

Is GigaCage Effective?

Current easiest method: make a fake JSArray...[\[ref\]](#)



Slightly limited R/W, but allows corrupting more complex structures elsewhere

Moving Towards A Heap Sandbox

Attackers will continue to find objects with corruptible pointers

Why not constrain the entire JS Heap?

- JavaScript manages many "external pointers" to browser memory

Moving Towards A Heap Sandbox

Attackers will continue to find objects with corruptible pointers

Why not constrain the entire JS Heap?

- JavaScript manages many "external pointers" to browser memory

Solution: Hold these pointers outside the heap and reference with index #

Future V8 Heap Sandbox

All JS objects confined to sandbox memory

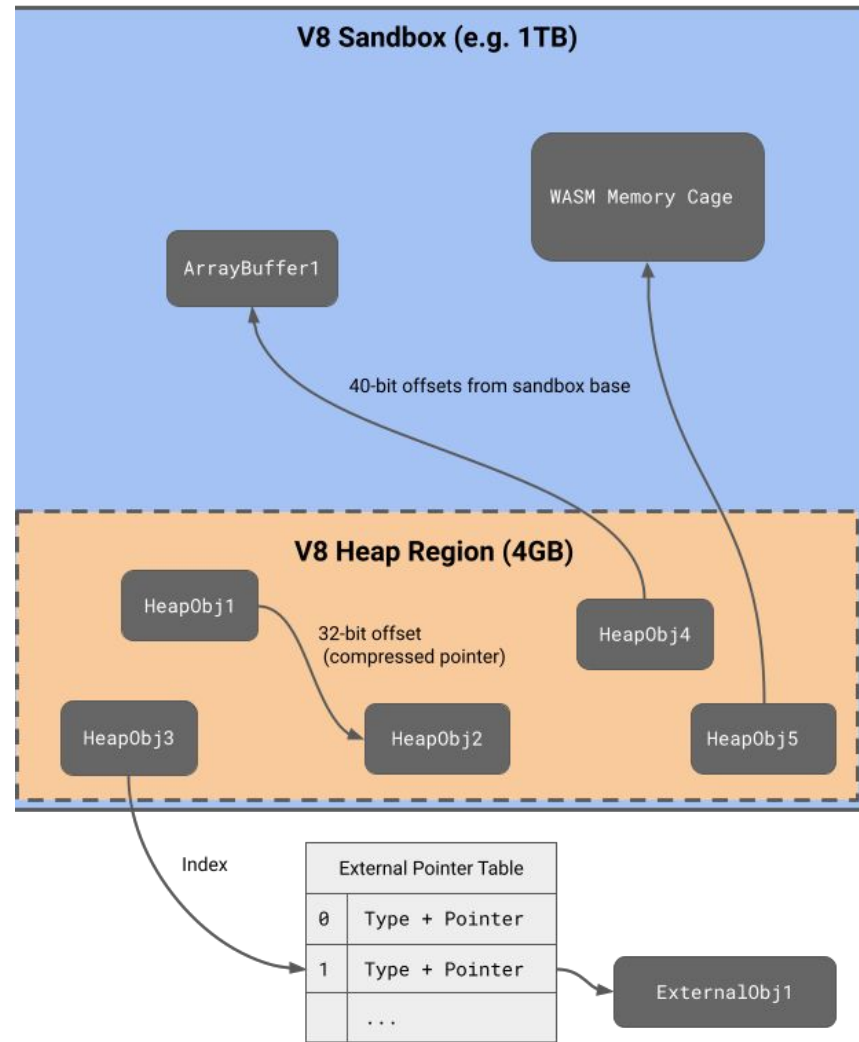
All other sensitive memory is outside:

- External pointers (and type) in table
- JIT compiler structures and code
- Any reference to other memory

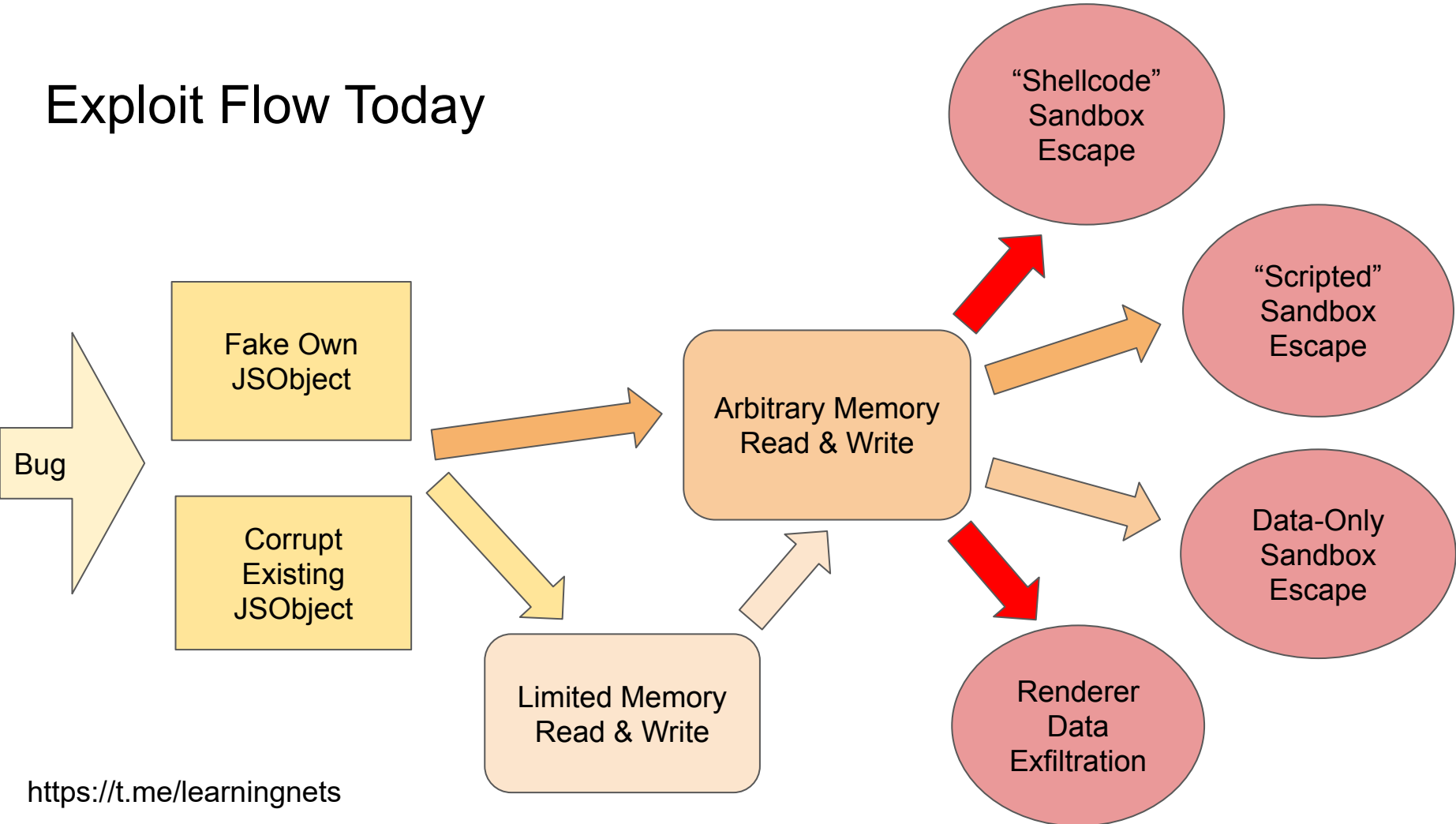
Exploit now relies on unsound behavior of external objects and code it has handles to

(similar to a sandbox escape...)

<https://t.me/learningnets>



Exploit Flow Today



What Have We Learned

Fewer bug classes, instead more “1-off” bugs, more complex JIT bugs

No significant changes to “early” exploitation phase (Same primitives available)

Current mitigations are not fully effective or applied evenly

Future mitigations seem more promising! (But still not bulletproof)