



Ring0 Crackme

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Keywords

Ring0



1. Abstract

This time, in new article, I'm not going to talk about any comercial protection, this time I will talk about one interesting crackme from www.crackmes.de which was submitted by Ms-Rem. This crackme has a lots of nice features that make it very very interesting, some of them are IDT/SDT hooks, rootkit tricks to hide process, usage of \Device\PhysicalMemory, 2 process execution, APC, etc...

In this tutorial I'll cover most of stuff handled by crackme and show you my steps for bypassing it. Also you will see some neet tricks when it comes to softice anti J , also this will maybe help you to understand how SoftICE can coexist with drivers that hook IDT. Anyway, before we start I wanna say that this crackme is real master piece, and big tnx to Ms-Rem for this crackme.

Tools that I'm using here are:

- SoftICE
- IDA
- LiveKd
- wARK
- LordPE
- Tasm32 β the best asm compiler ever

S verom u Boga, deroko/ARTeam



2. Know your enemy

Here is description of crackme published by Ms-Rem at www.crackmes.de:

This crackme use packing and ring0 tricks for hiding serial check code.

Crackme install driver dxapi32.sys.

Recommend to run this crackme on WMVare that will save you from problems with fall of the system.

So we already know the name of driver, starting crackme.exe for the first time will endup with MessageBox informing you that you have to reboot your system(driver is installed), so we restart system and after that crackme runs without a problem. Since we are dealing with driver we have to check our SDT and IDT just to be sure that those are ok before we start SoftICE.

SDT hooks:

```
KeServiceDescriptorTable           : 0x80552B80
KeServiceDescriptorTable.ServiceTable : 0x804DBD20
KeServiceDescriptorTable.ServiceLimit : 0x0000011C
```

```
SSDT entry 0019 hooked by : 0xF728AC9C - NtClose
SSDT entry 00BA hooked by : 0xF728AD70 - NtReadVirtualMemory
SSDT entry 0115 hooked by : 0xF728AD02 - NtWriteVirtualMemory
```

```
[!] Hooked Native Apis : 3
```

IDT hooks are only seen when crackme is loaded:

```
01h  0008:F728A7BC   Interrupt  32 bit    00    1
03h  0008:F728A78A   Interrupt  32 bit    03    1
```

Next thing is to analyze dxapi32.sys and get as much as possible informations about it. So we will move our efforts to new driver, entry point of driver seems normal, sets handlers for IRP_MJ_CREATE and IRP_MJ_WRITE, hooks SDT and creates hook check thread:

```
.text:00011744      call     native_api_services
.text:00011749      mov     edi, eax
.text:0001174B      test    edi, edi
.text:0001174D      jnz     short loc_11791
.text:0001174F      mov     esi, [ebp+DriverObject]
.text:00011752      push   esi                ; DriverObject
.text:00011753      call   create_device
.text:00011758      call   HookNativeApis
.text:0001175D      push   edi                ; StartContext
.text:0001175E      push   offset StartRoutine ; StartRoutine
.text:00011763      push   edi                ; ClientId
.text:00011764      push   edi                ; ProcessHandle
.text:00011765      push   edi                ; ObjectAttributes
.text:00011766      push   1F03FFh           ; DesiredAccess
.text:0001176B      lea   eax, [ebp+ThreadHandle]
.text:0001176E      push   eax                ; ThreadHandle
.text:0001176F      call   ds:PsCreateSystemThread

...

.text:00011784      mov     eax, offset CreateWriteHandler
.text:00011789      mov     [esi+48h], eax    ; IRP_MJ_WRITE
.text:0001178C      mov     [esi+38h], eax    ; IRP_MJ_CREATA
```



After that driver entrypoint is finished. We will examine most important parts of this initialization process.

```
.text:00010DDE native_api_services proc near
.text:00010DDE         mov     ecx, ds:NtBuildNumber
.text:00010DE4         movzx  ecx, word ptr [ecx]
.text:00010DE7         xor     eax, eax
.text:00010DE9         cmp    ecx, 893h
.text:00010DEF         jz     short win_2k
.text:00010DF1         cmp    ecx, 0A28h
.text:00010DF7         jz     short win_xp
.text:00010DF9         mov    eax, 0C0000002h
.text:00010DFE         retn
.text:00010DFF
.text:00010DFF win_xp:
.text:00010DFF         mov    dword_11B24, 7Ah
.text:00010E09         mov    NtClose_ServiceNum, 19h
.text:00010E13         mov    NtWVM_ServiceNumber, 115h
.text:00010E1D         mov    NtRVM_ServiceNumber, 0BAh
.text:00010E27         mov    dword_11B20, 84h
.text:00010E31         mov    dword_11ACC, 4
.text:00010E3B         mov    NtTerminateProcess_d, 101h
.text:00010E45         mov    dword_11ADC, 88h
.text:00010E4F         retn
.text:00010E50
.text:00010E50 win_2k
.text:00010E50         mov    dword_11B24, 6Ah
.text:00010E5A         mov    NtClose_ServiceNum, 18h
.text:00010E64         mov    NtWVM_ServiceNumber, 0F0h
.text:00010E6E         mov    NtRVM_ServiceNumber, 0A4h
.text:00010E78         mov    dword_11B20, 9Ch
.text:00010E82         mov    dword_11ACC, 8
.text:00010E8C         mov    NtTerminateProcess_d, 0E0h
.text:00010E96         mov    dword_11ADC, 0A0h
.text:00010EA0         retn
.text:00010EA0 native_api_services endp
```

Depending on Build number driver will retrieve service numbers which will be used as index into KeServiceDescriptorTable to hook APIs. You may easily determine which is which native API by disassembling ntdll.dll and watching which number is passed in eax before sysenter is used or int 2eh on win2k systems. (or you may write program to do this for you, it is very simple, check ntapasm supplied with this document), now we will examine hooking procedure of SDT:

```
.text:00010F24 HookNativeApis proc near
.text:00010F24
.text:00010F24 old_cr0 = dword ptr -4
.text:00010F24
.text:00010F24         push   ecx
.text:00010F25         push   esi
.text:00010F26         push   edi
.text:00010F27         cli
.text:00010F28         mov    eax, cr0
.text:00010F2B         mov    [esp+0Ch+old_cr0], eax
.text:00010F2F         and    eax, 0FFFFFFFh
.text:00010F34         mov    cr0, eax
.text:00010F37         mov    ecx, ds:KeServiceDescriptorTable
.text:00010F3D         mov    edx, [ecx]
.text:00010F3F         mov    eax, NtClose_ServiceNum
.text:00010F44         shl   eax, 2
.text:00010F47         mov    edx, [eax+edx]
.text:00010F4A         mov    Old_NtClose, edx
.text:00010F50         mov    ecx, [ecx]
.text:00010F52         mov    dword ptr [eax+ecx], offset NtCloseHook
.text:00010F59         mov    ecx, ds:KeServiceDescriptorTable
.text:00010F5F         mov    edx, [ecx]
.text:00010F61         mov    eax, NtWVM_ServiceNumber
.text:00010F66         shl   eax, 2
```



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```
.text:00010F69      mov     edx, [eax+edx]
.text:00010F6C      mov     Old_NtWriteVirtualMemory, edx
.text:00010F72      mov     ecx, [ecx]
.text:00010F74      mov     dword ptr [eax+ecx], offset hook_NtWriteVirtualMemory
.text:00010F7B      mov     ecx, ds:KeServiceDescriptorTable
.text:00010F81      mov     eax, NtRVM_ServiceNumber
.text:00010F86      mov     edx, [ecx]
.text:00010F88      shl     eax, 2
.text:00010F8B      mov     edx, [eax+edx]
.text:00010F8E      mov     Old_NtReadVirtualMemory, edx
.text:00010F94      mov     ecx, [ecx]
.text:00010F96      push   0 ; Remove
.text:00010F98      push   offset NotifyRoutine ; NotifyRoutine
.text:00010F9D      mov     dword ptr [eax+ecx], offset Hook_NrReadVirtualMemory
.text:00010FA4      call   PsSetCreateProcessNotifyRoutine
.text:00010FA9      mov     eax, ds:KeServiceDescriptorTable
.text:00010FAE      mov     eax, [eax+8] ; number of native APIS
.text:00010FB1      shl     eax, 2 ; multiply by 4
.text:00010FB4      push   206B6444h ; Tag
.text:00010FB9      push   eax ; NumberOfBytes
.text:00010FBA      push   0 ; PoolType
.text:00010FBC      mov     size_native_apis, eax
.text:00010FC1      call   ds:ExAllocatePoolWithTag
.text:00010FC7      mov     ecx, size_native_apis
.text:00010FCD      mov     edx, ds:KeServiceDescriptorTable
.text:00010FD3      mov     native_apis_pool, eax
.text:00010FD8      mov     esi, [edx] ; array of native APIS
.text:00010FDA      mov     edi, eax ; allocated pool
.text:00010FDC      mov     eax, ecx
.text:00010FDE      shr     ecx, 2
.text:00010FE1      rep movsd ; copy array of native APIS to allocated pool
.text:00010FE3      mov     ecx, eax ; this also copies hooked Native APIS
.text:00010FE5      and     ecx, 3
.text:00010FE8      rep movsb
.text:00010FEA      mov     eax, [esp+0Ch+old_cr0] ; restore WP in cr0
.text:00010FEE      mov     cr0, eax
.text:00010FF1      sti
.text:00010FF2      pop     edi
.text:00010FF3      pop     esi
.text:00010FF4      pop     ecx
.text:00010FF5      retn
.text:00010FF5 HookNativeApis endp
```

As you may see by looking at this code, crackme uses cr0 to wipe Write protection. Then it will hook NtWriteVirtualMemory, NtReadVirtualMemory and NtClose, also it will use PsSetCreateProcessNotifyRoutine to define callback that will be called whenever process in system is created/deleted. (process is object, so objects are created and deleted in windows J). After that whole list of API pointers from KeServiceDescriptorTable including hooked ones is copied to new nonpaged pool allocated before copying occurs. This copying occurs so SDT can't be restored easily, well not without a little bit of patching. Remember this PsSetCreateProcessNotifyRoutine, because you will understand why it is used later on... Also note two variables : native_apis_pool and size_native_apis where native_apis_pool is buffer with hooked SDT and size_native_apis is size of this pool (number of pointers * 4 obtained from KeServiceDescriptorTable):

```
typedef struct ServiceDescriptorEntry {
    unsigned int *ServiceTableBase;
    unsigned int *ServiceCounterTableBase;
    unsigned int NumberOfServices;
    unsigned char *ParamTableBase;
} SSDT_Entry;
```

Hooks of NtWrite/ReadVirtualMemory are used to deny write/read from/to protected process but NtClose is used as checker if driver is running or not:



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```
.text:00010C9C NtCloseHook:
.text:00010C9C      call    PsGetCurrentProcessId
.text:00010CA1      push   eax
.text:00010CA2      call   check_runing_process
.text:00010CA7      xor    ecx, ecx
.text:00010CA9      cmp    dword ptr [esp+4], 0FE12F2E5h
.text:00010CB1      setz   cl
.text:00010CB4      test   eax, ecx
.text:00010CB6      jz     short loc_10CC0
.text:00010CB8      mov    eax, 0EFFF2D7Fh
.text:00010CBD      retn   4
.text:00010CC0 loc_10CC0:
.text:00010CC0      jmp    Old_NtClose
```

Whenever NtClose is called with argument **0FE12F2E5h** driver will return **0EFFF2D7Fh** , but if driver isn't present then STATUS_INVALID_HANDLE will be returned, so this is used as check if driver is running J

Next what we are going to analyze is Thread created by driver startup code:

```
.text:00011596 ; void __stdcall StartRoutine(PVOID)
.text:00011596 StartRoutine  proc near
.text:00011596
.text:00011596 Object          = _KTIMER ptr -3Ch
.text:00011596 var_14         = qword ptr -14h
.text:00011596 var_C          = dword ptr -0Ch
.text:00011596 var_8          = dword ptr -8
.text:00011596 var_4          = dword ptr -4
.text:00011596
.text:00011596      push   ebp
.text:00011597      mov    ebp, esp
.text:00011599      sub    esp, 3Ch
.text:0001159C      push   ebx
.text:0001159D      push   1 ; Type
.text:0001159F      lea   eax, [ebp+Object]
.text:000115A2      push   eax ; Timer
.text:000115A3      call   ds:KeInitializeTimerEx
.text:000115A9      xor    ebx, ebx
.text:000115AB      push   ebx ; Dpc
.text:000115AC      push   3E8h ; Period
.text:000115B1      xor    eax, eax
.text:000115B3      push   eax
.text:000115B4      xor    ecx, ecx
.text:000115B6      push   ecx ; DueTime
.text:000115B7      lea   eax, [ebp+Object]
.text:000115BA      push   eax ; Timer
.text:000115BB      call   ds:KeSetTimerEx
.text:000115C1      cmp    byte_11AC0, bl
.text:000115C7      jnz   loc_11696
.text:000115CD      push   esi
.text:000115CE      push   edi
.text:000115CF      jmp    short loc_115D3
.text:000115D1 ; -----
.text:000115D1
.text:000115D1 loc_115D1:
.text:000115D1      xor    ebx, ebx
.text:000115D3
.text:000115D3 loc_115D3:
.text:000115D3      push   ebx ; Timeout
.text:000115D4      push   ebx ; Alertable
.text:000115D5      push   ebx ; WaitMode
.text:000115D6      push   ebx ; WaitReason
.text:000115D7      lea   eax, [ebp+Object]
.text:000115DA      push   eax ; Object
.text:000115DB      call   ds:KeWaitForSingleObject
.text:000115E1      cli
.text:000115E2      mov    eax, cr0
.text:000115E5      mov    [ebp+var_4], eax
.text:000115E8      and    eax, 0FFFFFFFh
.text:000115ED      mov    cr0, eax
.text:000115F0      mov    eax, ds:KeServiceDescriptorTable
.text:000115F5      mov    ecx, size_native_apis
```



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```
.text:000115FB      mov     edi, [eax]
.text:000115FD      mov     esi, native_apis_pool
.text:00011603      mov     eax, ecx
.text:00011605      shr     ecx, 2
.text:00011608      rep movsd
.text:0001160A      mov     ecx, eax
.text:0001160C      and     ecx, 3
.text:0001160F      rep movsb
.text:00011611      mov     eax, [ebp+var_4]
.text:00011614      mov     cr0, eax
.text:00011617      sti
.text:00011618      sidt   [ebp+var_14]
.text:0001161C      mov     ebx, dword ptr [ebp+var_14+2]
.text:0001161F      mov     si, [ebx+18h]
.text:00011623      rol     esi, 10h
.text:00011626      mov     si, [ebx+1Eh]
.text:0001162A      ror     esi, 10h
.text:0001162D      mov     [ebp+var_C], esi
.text:00011630      mov     eax, cr4
.text:00011633      and     eax, 0FFFFFFF7h
.text:00011636      mov     cr4, eax
.text:00011639      mov     si, [ebx+8]
.text:0001163D      rol     esi, 10h
.text:00011640      mov     si, [ebx+0Eh]
.text:00011644      ror     esi, 10h
.text:00011647      mov     [ebp+var_8], esi
.text:0001164A      xor     eax, eax
.text:0001164C      cmp     [ebp+var_8], offset new_int1_handle
.text:00011653      setnz  al
.text:00011656      xor     ecx, ecx
.text:00011658      cmp     [ebp+var_C], offset new_int3_handle
.text:0001165F      setnz  cl
.text:00011662      xor     ebx, ebx
.text:00011664      or      eax, ecx
.text:00011666      xor     ecx, ecx
.text:00011668      cmp     runin_processes_pid_ptr, ebx
.text:0001166E      setnz  cl
.text:00011671      test   eax, ecx
.text:00011673      jz     short loc_11688
.text:00011675      call  ds:IoGetCurrentProcess
.text:0001167B      push  10h          ; size_t
.text:0001167D      push  ebx          ; int
.text:0001167E      push  eax          ; void *
.text:0001167F      call  ds:memchr
.text:00011685      add     esp, 0Ch
.text:00011688      loc_11688:
.text:00011688      cmp     byte_11AC0, bl
.text:0001168E      jz     loc_115D1
.text:00011694      pop     edi
.text:00011695      pop     esi
.text:00011696      loc_11696:
.text:00011696      pop     ebx
.text:00011697      leave
.text:00011698      retn   4
.text:00011698      StartRoutine  endp
```

Here you may see how timer is set, but also note how, when certain amount of time passes it will overwrite SDT with saved copy of hooked SDT. Very nice, so if you restore SDT it will again hook it, and you will have to patch a few places to be able to dump parts of crackme memory. Patching isn't hard here, write driver that will nook rep movsd/rep movsb here:

```
.text:00011608      rep movsd
.text:0001160A      mov     ecx, eax
.text:0001160C      and     ecx, 3
.text:0001160F      rep movsb
```

You might wanna check anti-msrem.asm driver for complete source (not very advanced driver, only simple patching J).



After patching of this 2 instructions crackme won't run anymore, it will start but due to check with NtClose it won't run.

We also know that there are hooks of int1 and int3 so we may locate them in driver by using this simple formula:

$\text{Int1_address} - \text{driverbase} = 0000078\text{Ah}$ now go to that RVA in IDA and you will find int1 and int3 handles:

```
.text:0001078A new_int3_handle proc near
.text:0001078A         pusha
.text:0001078B         mov     di, 30h
.text:0001078F         mov     fs, di
.text:00010792         mov     eax, [esp+20h] ; EIP
.text:00010796         dec     eax             ; address of int3h in process
.text:00010797         push   eax
.text:00010798         call   handle_int3h
.text:0001079D         test   eax, eax
.text:0001079F         jnz    short __old_int3_handle
.text:000107A1         mov     di, 3Bh
.text:000107A5         mov     fs, di
.text:000107A8         popa
.text:000107A9         dec     byte ptr [esp+0] ; eip points to int3h
.text:000107AC         iret
.text:000107AD __old_int3_handle:
.text:000107AD         mov     di, 3Bh
.text:000107B1         mov     fs, di
.text:000107B4         popa
.text:000107B5         jmp     old_int3_handle
.text:000107B5 new_int3_handle endp
```

And int1:

```
.text:000107BC new_int1_handle proc near
.text:000107BC         arg_4   = dword ptr 8
.text:000107BC         pusha
.text:000107BD         push   fs
.text:000107BF         mov     di, 30h
.text:000107C3         mov     fs, di
.text:000107C6         push   eax
.text:000107C7         call   check_int1_process
.text:000107CC         test   eax, eax
.text:000107CE         jz     short __old_int1_handle
.text:000107D0         mov     eax, [esp+2Ch] ; eflags
.text:000107D4         and     eax, 0FFFFFFFh ; wipe TRAP flag if used
.text:000107D9         or     eax, 10000h     ; set Resume Flag
.text:000107DE         mov     [esp+2Ch], eax ; set eflags
.text:000107E2         xor     eax, eax
.text:000107E4         mov     dr7, eax      ; wipe dr7
.text:000107E7         mov     dr0, eax     ; wipe dr0
.text:000107EA         pop     fs
.text:000107EC         popa
.text:000107ED         iret
.text:000107EE __old_int1_handle:
.text:000107EE         pop     fs
.text:000107F0         popa
.text:000107F1         jmp     old_int1_handle
.text:000107F1 new_int1_handle endp
```

Int1 handle is not complicated to understand it will wipe dr7 and dr0 and Trap Flag from eflags if int1 handler is called from crackme, otherwise it will call old handler. Not much to talk about it, really simple.

On other hand int3 is much much more interesting because it has APC delivery in it. Really kewl.



In int3 handler there is procedure which I named handle_int3:

```
.text:000106CE handle_int3h   proc near
.text:000106CE int3_eip       = dword ptr  8
.text:000106CE
.text:000106CE          push     esi
.text:000106CF          call    PsGetCurrentProcessId
.text:000106D4          mov     esi, eax
.text:000106D6          push     esi
.text:000106D7          call    check_runing_process
.text:000106DC          test    eax, eax
.text:000106DE          jnz    short __good_process
.text:000106E0          inc     eax
.text:000106E1          jmp     short __not_my_process_0
.text:000106E3
.text:000106E3 __good_process:
.text:000106E3          push    ebx
.text:000106E4          push    esi
.text:000106E5          call    __check_runnin_processes
.text:000106EA          mov     esi, eax
.text:000106EC          xor     ebx, ebx
.text:000106EE          cmp     esi, ebx
.text:000106F0          jz     short __not_my_conditions
.text:000106F2          push    edi
.text:000106F3          push    206B6444h      ; Tag
.text:000106F8          push    10h           ; NumberOfBytes
.text:000106FA          push    ebx           ; PoolType
.text:000106FB          call    ds:ExAllocatePoolWithTag
.text:00010701          push    ebx           ; State
.text:00010702          mov     edi, eax      ; KEVENT type
.text:00010704          push    ebx           ; Type
.text:00010705          push    edi           ; Event
.text:00010706          call    ds:KeInitializeEvent
.text:0001070C          push    ebx           ; zero...
.text:0001070D          push    [esp+0Ch+int3_eip] ; address of int3h
.text:00010711          push    edi           ; PKEVENT
.text:00010712          push    [esi+proc_struct.user_apc] ; user procedure to be executed
.text:00010715          push    [esi+proc_struct.ptr_ethread] ; pointer to KTHREAD
.text:00010718          call    insert_apc    ; APC initialization routine = 1315474ch
.text:0001071D          push    ebx           ; Timeout
.text:0001071E          push    ebx           ; Alertable
.text:0001071F          push    ebx           ; WaitMode
.text:00010720          push    ebx           ; WaitReason
.text:00010721          push    edi           ; Object
.text:00010722          call    ds:KeWaitForSingleObject
.text:00010728          push    ebx           ; Tag
.text:00010729          push    edi           ; P
.text:0001072A          call    ds:ExFreePoolWithTag
.text:00010730          pop     edi
.text:00010731          jmp     short loc_10739
.text:00010733 __not_my_conditions:
.text:00010733          mov     large ds:0, ebx
.text:00010739
.text:00010739 loc_10739:
.text:00010739          xor     eax, eax
.text:0001073B          pop     ebx
.text:0001073C
.text:0001073C __not_my_process_0:
.text:0001073C          pop     esi
.text:0001073D          retn   4
.text:0001073D handle_int3h   endp
```

And insert_apc:

```
.text:00010626 insert_apc   proc near
.text:00010626 pkthread    = dword ptr  8
.text:00010626 user_procedure = dword ptr  0Ch
.text:00010626 EVENT_ptr    = dword ptr  10h
.text:00010626 int3h_eip    = dword ptr  14h
.text:00010626 set_to_zero  = dword ptr  18h
.text:00010626
.text:00010626          push    ebp
.text:00010627          mov     ebp, esp
```



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```
.text:00010629      push     edi
.text:0001062A      push     206B6444h      ; Tag
.text:0001062F      push     30h           ; NumberOfBytes
.text:00010631      push     0             ; PoolType
.text:00010633      call    ds:ExAllocatePoolWithTag
.text:00010639      push     [ebp+EVENT_ptr] ; IN PVOID NormalContext
.text:0001063C      mov     edi, eax
.text:0001063E      push     1             ; KPROCESSOR_MODE ApcMode (kernel user)
.text:00010640      push     [ebp+user_procedure] ; user APC routine
.text:00010643      push     0             ; RundownRoutine
.text:00010645      push     offset kernel_apc_routine ; kernel routine
.text:0001064A      push     0             ; KAPC_ENVIRONMENT
.text:0001064C      push     [ebp+pkthread] ; PKTHREAD
.text:0001064F      push     edi           ; P APC
.text:00010650      call    KeInitializeApc
.text:00010655      push     0             ; mode
.text:00010657      push     [ebp+set_to_zero] ; systemparam2
.text:0001065A      push     [ebp+int3h_eip] ; systemparam1
.text:0001065D      push     edi           ; P APC
.text:0001065E      call    KeInsertQueueApc
.text:00010663      pop     edi
.text:00010664      pop     ebp
.text:00010665      retn    14h
.text:00010665      insert_apc  endp ; sp = -30h
```

To get as much as possible information what is what I used LiveKd to dump live memory and to analyze addresses, that's how I found user APC procedure first. But later I saw that it could be located much easier (altho this is easy enough J)

To recognize processes (2 of them) driver is using one struct which I named proc_struct:

```
00000000 proc_struct      struc ; (sizeof=0x24)
00000000 Next           dd ?
00000004 unknown_0      dd ?
00000008 hidden_pid     dd ? ; pid of dialog process
0000000C apc_pid        dd ? ; pid of sleeping process
00000010 unknown_1      dd ?
00000014 ptr_eprocess     dd ? ; e process of hidden_pid
00000018 ptr_ethread      dd ? ; e thread of APC process
0000001C user_apc        dd ? ; user APC address
00000024 proc_struct     ends
```

Reason why EPROCESS is stored here is because process is hidden. And when process is hidden there is no other way to find it's eprocess, well at least PsLookupProcessByProcessId won't find it. But hidden processes can be located via SwapContext because windows is using thread scheduling, so yep hidden process will run as long as it's threads are visible J The reason why EPROCESS is stored is because crackme adopts ring0 WriteProcessMemory via KeAttachProcess, and without saved EPROCESS there will be no chance for crackme to connect to target process. APC process (I named it like that) is 2nd process of crackme which will loop with (from virtual.dll):

```
CODE:1315698E loc_1315698E:
CODE:1315698E      push     0FFFFFFFFh      ; bAlertable
CODE:13156990      push     0FFFFFFFFh      ; dwMilliseconds
CODE:13156992      call    SleepEx
CODE:13156997      jmp     short loc_1315698E
```

APC will be executed and SleepEx will return, but again it will fall into SleepEx loop, now you may also realize why there is PsSetCreateProcessNotifyRoutine. Simple, this 2nd process will never end, so you have to terminate it somehow and that's what Notify routine is doing. Here is small explanation from DDK:



PsSetCreateProcessNotifyRoutine adds a driver-supplied callback routine to, or removes it from, a list of routines to be called whenever a process is created or deleted.

NTSTATUS

```
PsSetCreateProcessNotifyRoutine(  
    IN PCREATE_PROCESS_NOTIFY_ROUTINE NotifyRoutine,  
    IN BOOLEAN Remove  
);
```

And prototype of NotifyRoutine is:

```
VOID  
(*PCREATE_PROCESS_NOTIFY_ROUTINE) (  
    IN HANDLE ParentId,  
    IN HANDLE ProcessId,  
    IN BOOLEAN Create  
);
```

The *ParentId* and *ProcessId* parameters identify the process, and the *Create* parameter indicates whether the process was created (TRUE) or deleted (FALSE).

So now driver knows when process is terminated and can terminate sleeping APC child. Lets see that routine (you may locate it by looking at hook_native_apis following references):

```
.text:00010EA2 ; void __stdcall NotifyRoutine(HANDLE,HANDLE,BOOLEAN)  
.text:00010EA2 NotifyRoutine  proc near  
.text:00010EA2  
.text:00010EA2 ObjectAttributes= OBJECT_ATTRIBUTES ptr -20h  
.text:00010EA2 CLIENT_ID      = CLIENT_ID ptr -8  
.text:00010EA2 Process_ID    = dword ptr  0Ch  
.text:00010EA2 create       = dword ptr  10h  
.text:00010EA2  
.text:00010EA2      push    ebp  
.text:00010EA3      mov     ebp, esp  
.text:00010EA5      sub     esp, 20h  
.text:00010EA8      push    ebx  
.text:00010EA9      xor     ebx, ebx  
.text:00010EAB      cmp     byte ptr [ebp+create], bl  
.text:00010EAE      jnz    short loc_10F1F  
.text:00010EB0      push    [ebp+Process_ID]  
.text:00010EB3      call   check_runing_process  
.text:00010EB8      test   eax, eax  
.text:00010EBA      jz     short loc_10F1F  
.text:00010EBC      push    [ebp+Process_ID]  
.text:00010EBF      call   __check_runnin_processes  
.text:00010EC4      cmp     eax, ebx  
.text:00010EC6      jz     short loc_10ECD  
.text:00010EC8      mov     ecx, [eax+proc_struct.apc_pid]  
.text:00010ECB      jmp    short loc_10ED8  
.text:00010ECD loc_10ECD:  
.text:00010ECD      push    [ebp+Process_ID]  
.text:00010ED0      call   find_proc_struct  
.text:00010ED5      mov     ecx, [eax+proc_struct.hiden_pid]  
.text:00010ED8  
.text:00010ED8 loc_10ED8:  
.text:00010ED8      push    eax  
.text:00010ED9      mov     [ebp+ObjectAttributes.Length], 18h  
.text:00010EE0      mov     [ebp+ObjectAttributes.RootDirectory], ebx  
.text:00010EE3      mov     [ebp+ObjectAttributes.Attributes], ebx  
.text:00010EE6      mov     [ebp+ObjectAttributes.ObjectName], ebx  
.text:00010EE9      mov     [ebp+ObjectAttributes.SecurityDescriptor], ebx  
.text:00010EEC      mov     [ebp+ObjectAttributes.SecurityQualityOfService], ebx  
.text:00010EEF      mov     [ebp+CLIENT_ID.UniqueProcess], ecx  
.text:00010EF2      mov     [ebp+CLIENT_ID.UniqueThread], ebx
```



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```
.text:00010EF5      call     sub_109DC
.text:00010EFA      lea     eax, [ebp+CLIENT_ID]
.text:00010EFD      push   eax                ; ClientId
.text:00010EFE      lea     eax, [ebp+ObjectAttributes]
.text:00010F01      push   eax                ; ObjectAttributes
.text:00010F02      push   1F0FFFh           ; DesiredAccess
.text:00010F07      lea     eax, [ebp+create]
.text:00010F0A      push   eax                ; ProcessHandle
.text:00010F0B      call   ds:ZwOpenProcess
.text:00010F11      test   eax, eax
.text:00010F13      jnz    short loc_10F1F
.text:00010F15      push   ebx                ; ExitStatus
.text:00010F16      push   [ebp+create]       ; ProcessHandle
.text:00010F19      call   ds:ZwTerminateProcess
.text:00010F1F      loc_10F1F:
.text:00010F1F      pop     ebx
.text:00010F20      leave
.text:00010F21      retn    0Ch
.text:00010F21 NotifyRoutine endp
```

What is going here? There are 3 procedures which have similar names `__check_runing_process`, `find_proc_struct` etc... common for them is that those are used to locate `proc_struct` for certain process by PID. I have presented partially reversed `proc_struct` so all that is done here is something like this:

First it checks if process is terminated, if so then it checks if terminated processes is protected process, if one of them is protected and is deleted then `NtTerminateProcess` is used to terminate other process. This will happen when you terminate child process (visible one from Task Manager or Process Explorer) or you click on 'X' in dialog box when Notify routine will kill Sleeping child process.

Now we have to examine `CreateWrite` handler in driver:

```
.text:00011414 CreateWriteHandler proc near                ; DATA XREF: start+B60
.text:00011414
.text:00011414 var_5C          = dword ptr -5Ch
.text:00011414 user_buffer_local = dword ptr -58h
.text:00011414 ms_exc         = CPPEH_RECORD ptr -18h
.text:00011414 Irp           = dword ptr 0Ch
.text:00011414
.text:00011414      push   4Ch
.text:00011416      push   offset unk_119B8
.text:0001141B      call   __SEH_prolog
.text:00011420      xor    ecx, ecx
.text:00011422      xor    esi, esi
.text:00011424      mov    ebx, [ebp+Irp]
.text:00011427      mov    eax, [ebx+60h] ; io_stack_location
.text:0001142A      mov    [ebx+IRP.IoStatus.Information], ecx
.text:0001142D      cmp    byte ptr [eax], 4 ; isl_MajorFunction
.text:00011430      jnz    short __create_request
.text:00011432      mov    eax, [eax+IO_STACK_LOCATION.Parameters.Write.Length]
.text:00011435      mov    esi, [ebx+IRP.UserBuffer]
.text:00011438      cmp    eax, 40h
.text:0001143B      jnz    short __wrong_write_size
.text:0001143D      mov    [ebp+ms_exc.disabled], ecx
.text:00011440      push   10h
.text:00011442      pop    ecx
.text:00011443      lea   edi, [ebp+user_buffer_local]
.text:00011446      rep movsd
.text:00011448      lea   eax, [ebp+user_buffer_local]
.text:0001144B      push  eax
.text:0001144C      call  main_user_request_handler
.text:00011451      mov    esi, eax
.text:00011453      mov    [ebp+var_5C], esi
.text:00011456      mov    [ebx+IRP.IoStatus.Information], 40h
.text:0001145D      loc_1145D:
.text:0001145D      or     dword ptr [ebp-4], 0FFFFFFFh
```



```
.text:00011461          jmp     short __create_request
.text:00011463
.text:00011463  loc_11463:
.text:00011463          xor     eax, eax
.text:00011465          inc     eax
.text:00011466          retn
.text:00011467  loc_11467:
.text:00011467          mov     esp, [ebp-18h]
.text:0001146A          mov     ebx, [ebp+Irp]
.text:0001146D          and     [ebx+IRP.IoStatus.Information], 0
.text:00011471          mov     esi, 0C0000006h
.text:00011476          jmp     short loc_1145D
.text:00011478  __wrong_write_size:
.text:00011478          mov     esi, 0C0000022h
.text:0001147D          mov     [ebx+IRP.IoStatus.Information], ecx
.text:00011480
.text:00011480  __create_request:
.text:00011480          mov     [ebx+IRP.IoStatus.anonymous_0.Status], esi
.text:00011483          xor     dl, dl          ; PriorityBoost
.text:00011485          mov     ecx, ebx       ; Irp
.text:00011487          call   ds:IoofCompleteRequest
.text:0001148D          mov     eax, esi
.text:0001148F          call   __SEH_epilog
.text:00011494          retn   8
.text:00011494  CreateWriteHandler endp
```

First CreateWriteHandler will check why is this driver called. Create or Write, if it is Write then you have copying of passed buffer to local variable and then that buffer is passed to main_user_request_handler:

```
.text:000112B6  main_user_request_handler proc
.text:000112B6
.text:000112B6  user_buffer      = dword ptr 8
.text:000112D3          mov     eax, [esp+user_buffer]
.text:000112D7          mov     ecx, [eax]
.text:000112D9          dec     ecx
.text:000112DA          cmp     ecx, 9          ; switch 10 cases
.text:000112DD          ja     short loc_1134B ; default
.text:000112DF          jmp     ds:off_11351[ecx*4] ; switch jump
```

From here you may see that buffer passed to driver via WriteFile is of maximum size 40h and 1st 4 bytes are used as Service ID, used later on with switch/case in main_user_request_handler.

Oki, we have basic knowledge of driver structure and how and what it is doing.

3. dumping virtual.exe and virtual.dll

If you trace crackme.exe and focs.dll a little bit you may notice how both of them are actually loaders for virtual.exe and virtual.dll, this extraction occurs before driver protection starts so you may dump regions from crackme. For this I'll go fast on this subject because it isn't very complex to dump those regions.

Load crackme.exe in debugger, BPX on VirtualAlloc, and watch how sections + pe header are being copied to it, after it is copied put process into infinite loop and dump region at 400000h, now all you have to do to get valid exe is to set RAW of section = RVA of section, now simple redirect process to ExitProcess and you have crackme dumped. (there is a lot more to fix!!!). Apply same trick to focs.dll and dump virtual.dll from it.



Good now you have virtual.exe and virtual.dll ready to be disassembled. First we try our dumped crackme.exe and it is stopped due to int3h in it, we know that there are int1/3 hooks so we also know that crackme will act on some specific way when int 3 occurs in it. It will deliver APC to child process and then it will wait to be signaled and return to process.

Now lets examine virtual.dll:

```
CODE:13157189      push    offset aDbgbreakpoint ; "DbgBreakPoint"
CODE:1315718E      push    offset aNtdll_dll ; "ntdll.dll"
CODE:13157193      call   GetModuleHandleA_0
CODE:13157198      push    eax
CODE:13157199      call   GetProcAddress_0
CODE:1315719E      mov     ds:DbgBreakPoint, eax
CODE:131571A3      mov     ds:byte_1315A994, 68h
CODE:131571AA      mov     ds:dword_1315A995, offset hook_DbgBreakPoint
CODE:131571B4      mov     ds:byte_1315A999, 0C3h
CODE:131571BB      push    offset NumberOfBytesWritten
CODE:131571C0      push    6
CODE:131571C2      push    offset byte_1315A994
CODE:131571C7      mov     eax, ds:DbgBreakPoint
CODE:131571CC      push    eax
CODE:131571CD      push    0FFFFFFFh
CODE:131571CF      call   WriteProcessMemory
CODE:131571D4      xor     eax, eax
CODE:131571D6      call   mess_with_callgate
CODE:131571DB      call   sub_13155520
CODE:131571E0      call   sub_131563FC
CODE:131571E5      call   install_driver
CODE:131571EA      push    offset NumberOfBytesWritten
CODE:131571EF      push    0
CODE:131571F1      push    0
CODE:131571F3      push    offset ThreadStart
CODE:131571F8      push    0
CODE:131571FA      push    0
CODE:131571FC      call   CreateThread
CODE:13157201      mov     ds:hThread, eax
```

Here you see concept of hooking some APIs in current process, interesting also procedures are mess_with_callgate which in this paticular case does nothing, only Sleep for 10h miliseconds, but in same procedure there is call to mess with gdt and to install callgate which isn't called.

After that we have install_driver, this procedure will perform search in system32\drivers\ folder looking for dxapi32.sys, if there is such file then crackme assumes that driver is installed and it will continue with execution, otherwise MessageBox will popup informing you that driver is installed and that you should reboot your system.

Next is CreateThread and new thread does nothing. It will call mess_with_callgate which will Sleep only and loop like that untill process isn't terminated.

Very interesting procedure is here:

```
CODE:131573BA      call   new_process
CODE:131573BF      test   al, al
CODE:131573C1      jz     short loc_131573CE
CODE:131573C3
CODE:131573C3  loc_131573C3:
CODE:131573C3      push    0FFFFFFFh
CODE:131573C5      push    0FFFFFFFh
CODE:131573C7      call   SleepEx
CODE:131573CC      jmp     short loc_131573C3
```

new_process procedure is responsible for creation of new process and also it will inform driver about parent/child pid. So lets see how it works.



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```
CODE:13156758 new_process   proc near           ; CODE XREF: CODE:131573BAp
CODE:13156758
CODE:13156758 var_7C             = dword ptr -7Ch
CODE:13156758 var_78             = dword ptr -78h
CODE:13156758 var_74             = dword ptr -74h
CODE:13156758 var_70             = dword ptr -70h
CODE:13156758 var_6C             = dword ptr -6Ch
CODE:13156758 var_68             = dword ptr -68h
CODE:13156758 Buffer             = dword ptr -64h
CODE:13156758 dwProcessId       = dword ptr -60h
CODE:13156758 hObject           = dword ptr -5Ch
CODE:13156758 hProcess          = dword ptr -58h
CODE:13156758 StartupInfo       = _STARTUPINFOA ptr -48h
CODE:13156758 NumberOfBytesWritten = dword ptr -4
CODE:13156758
CODE:13156758             push     ebp
CODE:13156759             mov     ebp, esp
CODE:1315675B             add     esp, 0FFFFFF84h
CODE:1315675E             push     ebx
CODE:1315675F             push     esi
CODE:13156760             push     edi
CODE:13156761             xor     eax, eax
CODE:13156763             mov     [ebp+var_74], eax
CODE:13156766             mov     [ebp+var_78], eax
CODE:13156769             mov     [ebp+var_7C], eax
CODE:1315676C             mov     [ebp+var_6C], eax
CODE:1315676F             mov     [ebp+var_70], eax
CODE:13156772             mov     [ebp+var_68], eax
CODE:13156775             xor     eax, eax
CODE:13156777             push     ebp
CODE:13156778             push     offset j_unknown_libname_10_72
CODE:1315677D             push     dword ptr fs:[eax]
CODE:13156780             mov     fs:[eax], esp
CODE:13156783             push     offset aZwvdmcontrol ; "ZwVdmControl"
CODE:13156788             push     offset aNtdll_dll_0 ; "ntdll.dll"
CODE:1315678D             call    GetModuleHandleA_0
CODE:13156792             push     eax ; hModule
CODE:13156793             call    GetProcAddress_0
CODE:13156798             mov     ebx, eax
CODE:1315679A             lea     eax, [ebp+NumberOfBytesWritten]
CODE:1315679D             push     eax ; lpNumberOfBytesRead
CODE:1315679E             push     0Ch ; nSize
CODE:131567A0             lea     eax, [ebp+Buffer]
CODE:131567A3             push     eax ; lpBuffer
CODE:131567A4             mov     edi, ebx
CODE:131567A6             push     edi ; lpBaseAddress
CODE:131567A7             push     0FFFFFFFFh ; hProcess
CODE:131567A9             call    ReadProcessMemory
CODE:131567AE             cmp     [ebp+Buffer], 64152917h <-- signature bytes
CODE:131567B5             jz     __2nd_process
```

Here we see that crackme checks if it is parent or child process by checking bytes in NtVdmControl, if no matching bytes were found, we are father process and it is time to create child process:

```
CODE:131567BB             lea     eax, [ebp+StartupInfo]
CODE:131567BE             mov     edx, 44h
CODE:131567C3             call    sub_13146004
CODE:131567C8             mov     [ebp+StartupInfo.cb], 44h
CODE:131567CF             lea     eax, [ebp+hProcess]
CODE:131567D2             push     eax ; lpProcessInformation
CODE:131567D3             lea     eax, [ebp+StartupInfo]
CODE:131567D6             push     eax ; lpStartupInfo
CODE:131567D7             push     0 ; lpCurrentDirectory
CODE:131567D9             push     0 ; lpEnvironment
CODE:131567DB             push     4 ; dwCreationFlags
CODE:131567DD             push     0 ; bInheritHandles
CODE:131567DF             push     0 ; lpThreadAttributes
CODE:131567E1             push     0 ; lpProcessAttributes
CODE:131567E3             lea     edx, [ebp+var_68]
```




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```
CODE:131567E6      xor     eax, eax
CODE:131567E8      call   @ParamStr
CODE:131567ED      mov     eax, [ebp+var_68]
CODE:131567F0      call   unknown_libname_21
CODE:131567F5      push   eax                ; lpCommandLine
CODE:131567F6      push   0                  ; lpApplicationName
CODE:131567F8      call   CreateProcessA
CODE:131567FD      mov     [ebp+Buffer], 64152917h <-- signature bytes
CODE:13156804      call   GetCurrentProcessId
CODE:13156809      mov     esi, eax
CODE:1315680B      mov     [ebp+dwProcessId], esi <-- buffer + 4 = parent PID
CODE:1315680E      lea    edx, [ebp+var_70]
CODE:13156811      mov     eax, esi
CODE:13156813      call   @Sysutils@IntToStr$qqri_0
CODE:13156818      mov     ecx, [ebp+var_70]
CODE:1315681B      lea    eax, [ebp+var_6C]
CODE:1315681E      mov     edx, offset aFocs ; "FOCS"
CODE:13156823      call   @System@LStrCat3$qqrv
CODE:13156828      mov     eax, [ebp+var_6C]
CODE:1315682B      call   unknown_libname_21
CODE:13156830      push   eax                ; lpName
CODE:13156831      push   0                  ; bInitialState
CODE:13156833      push   0FFFFFFFFh        ; bManualReset
CODE:13156835      push   0                  ; lpEventAttributes
CODE:13156837      call   CreateEventA
CODE:1315683C      mov     esi, eax
CODE:1315683E      push   2                  ; dwOptions
CODE:13156840      push   0                  ; bInheritHandle
CODE:13156842      push   0                  ; dwDesiredAccess
CODE:13156844      lea    eax, [ebp+hObject]
CODE:13156847      push   eax                ; lpTargetHandle
CODE:13156848      mov     [ebp+hProcess]
CODE:1315684B      push   eax                ; hTargetProcessHandle
CODE:1315684C      push   esi                ; hSourceHandle
CODE:1315684D      push   0FFFFFFFFh        ; hSourceProcessHandle
CODE:1315684F      call   DuplicateHandle
CODE:13156854      lea    eax, [ebp+NumberOfBytesWritten]
CODE:13156857      push   eax                ; lpNumberOfBytesWritten
CODE:13156858      push   0Ch                ; nSize
CODE:1315685A      lea    eax, [ebp+Buffer]
CODE:1315685D      push   eax                ; lpBuffer
CODE:1315685E      push   edi                ; lpBaseAddress
CODE:1315685F      mov     eax, [ebp+hProcess]
CODE:13156862      push   eax                ; hProcess
CODE:13156863      call   WriteProcessMemory
CODE:13156868      mov     eax, [ebp-54h]
CODE:1315686B      push   eax                ; hThread
CODE:1315686C      call   ResumeThread
CODE:13156871      push   0FFFFFFFFh        ; dwMilliseconds
CODE:13156873      push   esi                ; hHandle
CODE:13156874      call   WaitForSingleObject
CODE:13156879      push   esi                ; hObject
CODE:1315687A      call   CloseHandle
CODE:1315687F      xor     ebx, ebx
CODE:13156881      jmp    loc_13156999
```

As you may see father will tell child it's PID, and will continue with initialization, now observe how child is acting, it will collect PIDs of both processes and inform driver about them using service ID : 7, also it will send to driver address of APC procedure and will enter into infinite loop with SleepEx allowing in such way APC to "interrupt" it's sleeping state.

```
CODE:13156886
CODE:13156886      __2nd_process:
CODE:13156886      mov     eax, [ebp+dwProcessId]
CODE:13156889      push   eax                ; dwProcessId
CODE:1315688A      push   0                  ; bInheritHandle
CODE:1315688C      push   1F0FFFh           ; dwDesiredAccess
CODE:13156891      call   OpenProcess
CODE:13156896      mov     ds:pHandle, eax
CODE:1315689B      push   0                  ; hTemplateFile
CODE:1315689D      push   0                  ; dwFlagsAndAttributes
CODE:1315689F      push   3                  ; dwCreationDisposition
CODE:131568A1      push   0                  ; lpSecurityAttributes
CODE:131568A3      push   0                  ; dwShareMode
```




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```
CODE:131568A5      push      40000000h      ; dwDesiredAccess
CODE:131568AA      push      offset a_RootSystem0_4 ;"\\\\.\\Root#SYSTEM#0000#{"...
CODE:131568AF      call     CreateFileA
CODE:131568B4      mov      ds:hDevice, eax
CODE:131568B9      mov      eax, [ebp+dwProcessId]
CODE:131568BC      mov      ds:parent_pid, eax
CODE:131568C1      call     GetCurrentProcessId
CODE:131568C6      mov      ds:current_pid, eax
CODE:131568CB      mov      ds:Buffer, 7
CODE:131568D5      mov      eax, offset dword_1315A900
CODE:131568DA      mov      ds:dword_1315A8C4, eax
CODE:131568DF      mov      eax, offset APC_userprocedure
CODE:131568E4      mov      ds:apc_procedure_addr, eax
CODE:131568E9      push     0              ; lpOverlapped
CODE:131568EB      lea     eax, [ebp+NumberOfBytesWritten]
CODE:131568EE      push     eax            ; lpNumberOfBytesWritten
CODE:131568EF      push     40h           ; nNumberOfBytesToWrite
CODE:131568F1      push     offset Buffer   ; lpBuffer
CODE:131568F6      mov     eax, ds:hDevice
CODE:131568FB      push     eax            ; hFile
CODE:131568FC      call    WriteFile_0
CODE:13156901      push     0              ; hTemplateFile
CODE:13156903      push     0              ; dwFlagsAndAttributes
CODE:13156905      push     3              ; dwCreationDisposition
CODE:13156907      push     0              ; lpSecurityAttributes
CODE:13156909      push     0              ; dwShareMode
CODE:1315690B      push     80000000h     ; dwDesiredAccess
CODE:13156910      lea     edx, [ebp+var_7C]
CODE:13156913      xor     eax, eax
CODE:13156915      call    @ParamStr
CODE:1315691A      mov     eax, [ebp+var_7C]
CODE:1315691D      lea     edx, [ebp+var_78]
CODE:13156920      call    unknown_libname_89
CODE:13156925      mov     ecx, [ebp+var_78]
CODE:13156928      lea     eax, [ebp+var_74]
CODE:1315692B      mov     edx, offset dword_13156A34
CODE:13156930      call    @System@LStrCat3$qqrv
CODE:13156935      mov     eax, [ebp+var_74]
CODE:13156938      call    unknown_libname_21
CODE:1315693D      push     eax            ; lpFileName
CODE:1315693E      call    CreateFileA
CODE:13156943      mov     ebx, eax
CODE:13156945      cmp     ebx, 0FFFFFFFh
CODE:13156948      jnz     short loc_13156963
CODE:1315694A      mov     eax, [ebp+hObject]
CODE:1315694D      push     eax            ; hEvent
CODE:1315694E      call    SetEvent
CODE:13156953      mov     eax, [ebp+hObject]
CODE:13156956      push     eax            ; hObject
CODE:13156957      call    CloseHandle
CODE:1315695C      push     0FFFFFFFh     ; dwMilliseconds
CODE:1315695E      call    Sleep
CODE:13156963      loc_13156963:
CODE:13156963      mov     ecx, offset unk_1315A8BC
CODE:13156968      xor     edx, edx        ; lDistanceToMove
CODE:1315696A      mov     eax, ebx        ; hFile
CODE:1315696C      call    sub_131547D8
CODE:13156971      mov     ds:dword_1315A924, eax
CODE:13156976      push     ebx            ; hObject
CODE:13156977      call    CloseHandle
CODE:1315697C      mov     eax, [ebp+hObject]
CODE:1315697F      push     eax            ; hEvent
CODE:13156980      call    SetEvent
CODE:13156985      mov     eax, [ebp+hObject]
CODE:13156988      push     eax            ; hObject
CODE:13156989      call    CloseHandle
CODE:1315698E      loc_1315698E:
CODE:1315698E      push     0FFFFFFFh     ; bAlertable
CODE:13156990      push     0FFFFFFFh     ; dwMilliseconds
CODE:13156992      call    SleepEx
CODE:13156997      jmp     short loc_1315698E
CODE:13156999      ; -----
CODE:13156999      loc_13156999:
```



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```
CODE:13156999      xor     eax, eax
CODE:1315699B      pop     edx
CODE:1315699C      pop     ecx
CODE:1315699D      pop     ecx
CODE:1315699E      mov     fs:[eax], edx
CODE:131569A1      push   offset loc_131569BB
CODE:131569A6      loc_131569A6:
CODE:131569A6      lea    eax, [ebp+var_7C]
CODE:131569A9      mov     edx, 6
CODE:131569AE      call   unknown_libname_12
CODE:131569B3      retn
CODE:131569B3      new_procesan  endp / sp = -90h
```

So you have basic knowledge how crackme is working. I'll talk only about steps that I've taken to unpack this crackme, now I have realized that some steps weren't necessary but what the hack...

By looking at user APC procedure stored in virtual.dll I noticed this part of code which will call driver:

```
CODE:1315461C      attach_to_process_r0 proc near
CODE:1315461C      NumberOfBytesWritten= dword ptr -44h
CODE:1315461C      Buffer          = dword ptr -40h
CODE:1315461C      var_3C         = dword ptr -3Ch
CODE:1315461C      var_38         = dword ptr -38h
CODE:1315461C      var_34         = dword ptr -34h
CODE:1315461C
CODE:1315461C      add     esp, 0FFFFFFBCh
CODE:1315461F      mov     [esp+44h+Buffer], 0Ah
CODE:13154627      mov     [esp+8], edx
CODE:1315462B      mov     [esp+12], eax
CODE:1315462F      mov     [esp+10h], ecx
CODE:13154633      push   0          ; lpOverlapped
CODE:13154635      lea    eax, [esp+48h+NumberOfBytesWritten]
CODE:13154639      push   eax          ; lpNumberOfBytesWritten
CODE:1315463A      push   40h         ; nNumberOfBytesToWrite
CODE:1315463C      lea    eax, [esp+50h+Buffer]
CODE:13154640      push   eax          ; lpBuffer
CODE:13154641      mov     eax, ds:hDevice
CODE:13154646      push   eax          ; hFile
CODE:13154647      call   WriteFile_0
CODE:1315464C      add     esp, 44h
CODE:1315464F      retn
CODE:1315464F      attach_to_process_r0 endp
```

As you may see it is using Service ID 0Ah to tell driver something, and soon we will find out what, but we have to switch to driver disassembly now and take a look at switch/case in main_user_request_handler here:

```
.text:00011331      push   dword ptr [eax+0Ch] ; case 0x9
.text:00011334      push   dword ptr [eax+8] ; destination
.text:00011337      push   dword ptr [eax+4] ; source
.text:0001133A      call   PsGetCurrentProcessId
.text:0001133F      push   eax
.text:00011340      call   find_proc_struct
.text:00011345      push   eax          ; proc_struct
.text:00011346      call   copy_memory_from_r0
```

And we enter into copy_memory_from_r0:

```
.text:0001083C ; int __stdcall copy_memory_from_r0(int proc_struct,int source,int destination,SIZE_T
NumberOfBytes)
.text:0001083C      copy_memory_from_r0 proc near          .text:0001083C
.text:0001083C      proc_struct      = dword ptr 8
.text:0001083C      source           = dword ptr 0Ch
.text:0001083C      destination      = dword ptr 10h
```



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```
.text:0001083C NumberOfBytes = dword ptr 14h
.text:0001083C
.text:0001083C      push    ebp
.text:0001083D      mov     ebp, esp
.text:0001083F      push    ebx
.text:00010840      push    esi
.text:00010841      push    edi
.text:00010842      push    206B6444h          ; Tag
.text:00010847      push    [ebp+NumberOfBytes] ; NumberOfBytes
.text:0001084A      push    0                  ; PoolType
.text:0001084C      call   ds:ExAllocatePoolWithTag
.text:00010852      mov     ecx, [ebp+NumberOfBytes]
.text:00010855      mov     esi, [ebp+source]
.text:00010858      mov     ebx, eax
.text:0001085A      mov     eax, ecx
.text:0001085C      shr     ecx, 2
.text:0001085F      mov     edi, ebx          ; from esi to newly allocated nonpaged pool
.text:00010861      rep movsd
.text:00010863      mov     ecx, eax
.text:00010865      mov     eax, [ebp+proc_struct]
.text:00010868      and     ecx, 3
.text:0001086B      rep movsb
.text:0001086D      push    [eax+proc_struct.ptr_eprocess] ; e process of hidden_pid
.text:00010870      call   KeAttachProcess
.text:00010875      cli
.text:00010876      mov     eax, cr0
.text:00010879      mov     [ebp+source], eax ; save cr0
.text:0001087C      and     eax, 0FFFFFFFh ; wipe WP protection
.text:00010881      mov     cr0, eax
.text:00010884      mov     ecx, [ebp+NumberOfBytes]
.text:00010887      mov     edi, [ebp+destination]
.text:0001088A      mov     eax, ecx
.text:0001088C      shr     ecx, 2
.text:0001088F      mov     esi, ebx
.text:00010891      rep movsd
.text:00010893      mov     ecx, eax
.text:00010895      and     ecx, 3
.text:00010898      rep movsb
.text:0001089A      mov     eax, [ebp+source]
.text:0001089D      mov     cr0, eax
.text:000108A0      sti
.text:000108A1      call   KeDetachProcess
.text:000108A6      push    0                  ; Tag
.text:000108A8      push    ebx                ; P
.text:000108A9      call   ds:ExFreePoolWithTag
.text:000108AF      pop     edi
.text:000108B0      pop     esi
.text:000108B1      pop     ebx
.text:000108B2      pop     ebp
.text:000108B3      retn   10h
.text:000108B3      copy_memory_from_r0 endp ; sp = -4
```

copy_memory_from_r0 is custom implementation of WriteProcessMemory. So APC procedure once it locates code that should be copied to hidden process will tell driver what to copy and where. Here you see usage of KeAttachProcess, but maybe some can ask why it is first copying data to nonpaged pool instead of direct copy? Answer is simple, KeAttachProcess WILL change value of cr3 which has physical address of page directory entry. When it happens source address of 2nd process won't be visible to us, and copying from nonexisting address will result in page_fault, and remember what page_fault in ring0 means.

So that's it about static disassembly, now we are moving to our activities with SoftICE, it is much easier when you can check your theory in practice. But still SoftICE can't be loaded due to hooks in IDT. Now comes my workaround to make softice work with this crackme without any sideeffects on its execution.



4. SoftICE comes to rescue

I'm usually not impressed by my ideas but I really loved this one, and I was thinking about making SoftICE to work with this crackme for a 2-3 hours, finally I set down and wrote everything on paper and started writing my driver to enable presence of SoftICE.

In driver there is only one procedure responsible for interrupt hooking which is called to hook and unhook interrupts:

```
.text:000103FE hook_interrupt  proc near
.text:000103FE
.text:000103FE var_C          = qword ptr -0Ch
.text:000103FE var_4          = dword ptr -4
.text:000103FE int_vector     = dword ptr  8      <-- arg0
.text:000103FE int_handler    = dword ptr  0Ch     <-- arg1
.text:000103FE arg_8          = word ptr  10h     <-- arg2
...
.text:0001042D                cli
.text:0001042E                call   sub_1039C    <-- real hook occurs here
.text:00010433                sti
.text:00010434                mov    eax, [ebp+var_4] <-- old handle
.text:00010437                pop    edi
.text:00010438                pop    esi
.text:00010439                pop    ebx
.text:0001043A                leave
.text:0001043B                retn   0Ch
.text:0001043B hook_interrupt  endp
```

This procedure is relatively simple and is invoked from here, and also from Unhook_interrupts:

```
.text:000107F8 hookints_drs  proc near
.text:000107F8                push  0EEh
.text:000107FD                push  offset new_int3_handle
.text:00010802                push  3
.text:00010804                call  hook_interrupt
.text:00010809                push  8Eh
.text:0001080E                push  offset new_int1_handle
.text:00010813                push  1
.text:00010815                mov   old_int3_handle, eax
.text:0001081A                call  hook_interrupt
.text:0001081F                mov   old_int1_handle, eax
.text:00010824                cli
.text:00010825                mov   eax, cr4
.text:00010828                and   eax, 0FFFFFFF7h ; DE bit wiped out
.text:0001082B                mov   cr4, eax
.text:0001082E                xor   eax, eax
.text:00010830                mov   dr7, eax        ; dr7 to 0 (no local nor global bpms)
.text:00010833                mov   dr1, eax        ; dr1 and dr0 are also zeroed
.text:00010836                mov   dr0, eax
.text:00010839                sti
.text:0001083A                retn
.text:0001083A hookints_drs  endp
```

As you may see hook interrupt will return address of old handler, and hook it with passed one, so since interrupt hooking isn't nuclear physics we may detour this procedure to our driver, and we are also going to detour int handlers in dxapi32.sys so both of them will point to SoftICE handlers.

This is accomplished by detouring crackme int 1/3 handlers to handlers installed by SoftICE, on other hand SoftICE handler look like this:

```
push  _KiTrapXX
jmp   __softice
```



So if SoftICE doesn't handle int1/int3 or other trap/fault it will call KiTrapXX that is pushed on stack, now comes my idea. But this time we are going to use smart approach on hooking driver because one mistake and you will get BSOD (not from fact that driver wanna kill you but because SoftICE installs some int3h in some native APIs to keep track of something – exports + KeServiceDescriptorTable!? Wrongly handled int3h will BSOD system)...

Oki I will rip 5 bytes from dxapi32.sys int1/int3 handler and will detour them to SoftICE handler, but I will also change:

```
                push    __KiTrapXX            >>>>>>>>>>>> push    __stolen_bytesint1/3>>>>
                jmp     __softice            >>>>>>>>>>>> jmp     __softice                V
                                                V
__stolen_bytes1/3:    pusha   <<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<
                    mov     di, 30h
                    mov     fs, di
                    nop
                    ...
                    push    __bytes_after_hook
                    retn
```

Maybe disassembly from softice will make more sense?

```
:idt
Int  Type      Sel:Offset      Attributes Symbol/Owner
IDTbase=8003F400 Limit=07FF
0000 IntG32     0008:804D8BFF DPL=0 P   _KiTrap00
0001 IntG32     0008:F728A7BC DPL=3 P   dxapi32!.text+04BC
0002 TaskG      0058:00000000 DPL=0 P   _KiTrap02
0003 IntG32     0008:F728A78A DPL=3 P   dxapi32!.text+048A
0004 IntG32     0008:804D92E0 DPL=3 P   _KiTrap04
...
:idt 3
0003 IntG32     0008:F728A78A DPL=3 P   dxapi32!.text+048A
:u f728a78a
0008:F728A78A JMP     _F8AB14EC           <-- detour to SoftICE
0008:F728A78F MOV     FS,DI
0008:F728A792 MOV     EAX,[ESP+20]
0008:F728A796 DEC     EAX
0008:F728A797 PUSH    EAX
0008:F728A798 CALL   _F728A6CE
0008:F728A79D TEST   EAX,EAX
0008:F728A79F JNZ    _F728A7AD
...
:u f8ab14ec
0008:F8AB14EC PUSH    F8AB14BE           <-- 5bytes + jmp back to dxapi32.sys
0008:F8AB14F1 JMP     _F87B770D           <-- jmp __SoftICE
0008:F8AB14F6 MOV     EAX,[ESP+04]
0008:F8AB14FA MOV     EDI,[F8AB12A4]
0008:F8AB1500 LEA    EDI,[EAX*8+EDI]
0008:F8AB1503 MOVZX   EAX,WORD PTR [EDI+06]
0008:F8AB1507 SHL    EAX,10
0008:F8AB150A MOV     AX,[EDI]
...
:u f8ab14be
0008:F8AB14BE PUSHAD                <-- stolen bytes overwritten by
0008:F8AB14BF MOV     DI,0030         <-- hook to SoftICE
0008:F8AB14C3 NOP
```



```
0008:F8AB14C4  NOP
0008:F8AB14D7  NOP
0008:F8AB14D8  NOP                <-- BPX here (or any nop) if you
0008:F8AB14D9  NOP                wanna debug dxapi32.sys int3
0008:F8AB14DA  NOP                handler
0008:F8AB14DB  NOP
0008:F8AB14DC  PUSH             F728A78F        <-- jmp to rest of code
0008:F8AB14E1  RET
...
:u f728a78f
0008:F728A78F  MOV             FS,DI        <-- dxapi32.sys int3h handle
0008:F728A792  MOV             EAX,[ESP+20]
0008:F728A796  DEC             EAX
0008:F728A797  PUSH           EAX
0008:F728A798  CALL           _F728A6CE
0008:F728A79D  TEST           EAX,EAX
0008:F728A79F  JNZ           _F728A7AD
0008:F728A7A1  MOV            DI,003B
0008:F728A7A5  MOV            FS,DI
0008:F728A7A8  POPAD
0008:F728A7A9  DEC            BYTE PTR [ESP]
0008:F728A7AC  IRETD
0008:F728A7AD  MOV            DI,003B
0008:F728A7B1  MOV            FS,DI
0008:F728A7B4  POPAD
0008:F728A7B5  JMP            [_KiTrap03]    <-- KiTrap03 handle
```

Looks nice J So SoftICE handler will always execute first, giving chance to SoftICE to handle single stepping and step overs without a problem, now I can also debug int3 handle of dxapi32.sys, isn't it nice. Also I have to hook interrupt_hook so I can control when crackme wants to hook and to unhook interrupt.

If hooking is performed then write to SoftICE hook addresses of detoured handlers, and return address of KiTrap01 or KiTrap03 depending which int is being hooked, if unhooking is performed then write to SoftICE handler (push) KiTrap01 and KiTrap03 and in restore SoftICE interrupts in IDT. If we don't return KiTrap01/KiTrap03 during this hooking then there is chance of falling into deadlock...

Example -> driver hooks IDT and takes old handler from there -> driver handle is hooked to point to SoftICE -> int3 -> softICE handler takes over control and it doesn't handle it -> driver takes control and decides that this is not int 3h in hidden process so it will call old handle -> but wait!? Old handle is softice handle which doesn't handle int3h so it calls again dxapi32.sys handle and so on, so on...

For detailed implementation check anti-msrem1.asm driver, it works like charm, and thanks to it I could better undersand my target.



5. Removing int3h from crackme.exe

As I mentioned earlier, crackme will use APC in second process to handle int3h in hidden process. To get better understand of what is happening here I have used i3here on in softice to break when int3h occurs (remember SoftICE handles int3 first with anti-msrem1.asm driver loaded), next thing for me was to set bpx at copy_memory_from_r0 in Write handler and to observe what is going on. Maybe I could analyze better APC procedure but I really liked my idea about SoftICE presence and I've decided to give it a shot.

I have set breaks and watched what is copied, where and how.

What I saw was more than enough to get right conclusion. First crackme will overwrite all stolen code block with int3h, it is logical step because only int3h will force it's restoring, but after all blocks are overwritten with int3h, then real code will be copied over one of int3h blocks. This copying is used so we never get fully reconstructed code in crackme. Ms-Rem was very nice and has included real serial with this crackme so we can log all places where stolen codes are. Anyway, real serial wasn't necessary because by int3h buffer overwriting we could find all places where stolen code is. Without real key we could locate all buffers but then we should manually redirect execution to those places and get real code or else better analyze APC procedure in virtual.dll.

First to get all buffers I hooked copy_memory_from_r0 (check log.asm driver) and logged where are all addresses:

```
00000000    0.00000000 Destination - 0x00452325
00000001    0.00003297 Destination - 0x00452365
00000002    0.00005364 Destination - 0x0045242C
00000003    0.00007347 Destination - 0x00452486
00000004    0.00009303 Destination - 0x00452548
00000005    0.00011286 Destination - 0x004525BD
00000006    0.00013270 Destination - 0x00452684
00000007    0.00015253 Destination - 0x004526DF
00000008    0.00017237 Destination - 0x00452722
00000009    0.00019192 Destination - 0x0045276D
00000010    0.00021148 Destination - 0x004527B4
00000011    0.00023103 Destination - 0x0045281B
00000012    0.00025087 Destination - 0x004528BF
```

My hook was stored here:

```
.text:00010884          mov     ecx, [ebp+NumberOfBytes]
.text:00010887          mov     edi, [ebp+destination]
.text:0001088A          mov     eax, ecx    <-- hook here
.text:0001088C          shr     ecx, 2     <-- and here
.text:0001088F          mov     esi, ebx
.text:00010891          rep movsd
.text:00010893          mov     ecx, eax
.text:00010895          and     ecx, 3
.text:00010898          rep movsb
```

Here we have more than enough for hooking(5 bytes), we have size of copied block, and we have also destination buffer in EDI and source buffer in EBX, so now I can log all buffers and also dump them into files and later simply apply those dumps to dumped crackme and it is unpacked.

Again I will suggest to check my log.asm driver, it isn't very complex so you will be able to understand it easily. That's pretty much about dumping and fixing.



6. Process hiding from ring3

Even if target is fully deprotected, there are still some more stuff left to fix, one of them is to disable process hiding from ring3. I will sum it here, but for detaild explanation about this trick you should refer to crazylord's paper in Phrack named "Playing with windows /dev/(k)mem" [1].

First open your target in IDA and find reference to ZwOpenSection, you will see also how privileges for \Device\PhysicalMemory are being adjusted so it can be opened and modified, all of this is documented in paper I mentioned.

Then it enters into GDT modification and installs this callgate:

```
CODE:004518CE      mov     eax, offset call_gate_procedure
CODE:004518D3      mov     edx, [ebx]
CODE:004518D5      mov     [edx], ax
CODE:004518D8      shr     eax, 10h
CODE:004518DB      mov     edx, [ebx]
CODE:004518DD      mov     [edx+6], ax
CODE:004518E1      mov     eax, [ebx]
CODE:004518E3      mov     word ptr [eax+4], 0EC00h
CODE:004518E9      xor     eax, eax
CODE:004518EB      mov     dword ptr ds:callgate_descriptor, eax
CODE:004518F0      mov     ax, [ebx]
CODE:004518F3      sub     ax, word ptr ds:lpBaseAddress
CODE:004518FA      sub     ax, word ptr [ebp+var_C]
CODE:004518FE      mov     word ptr ds:callgate_descriptor+4, ax
```

As you may see in this code it writes address of call_gate_procedure to free slot in GDT and also stores Selector into callgate_descriptor+4. When we are calling callgate only thing that is relevant in it is Selector while address is not important. By executing far call to that callgate CPU will automaticaly redirect execution to address sotred in GDT.

Next thing we are going to examine is call_gate_procedure:

```
CODE:004516A4 call_gate_procedure:
CODE:004516A4      cli
CODE:004516A5      pusha
CODE:004516A6      pushf
CODE:004516A7      mov     di, 30h
CODE:004516AB      mov     fs, di
CODE:004516AE      call   ds:callgate_current
CODE:004516B4      mov     di, 3Bh
CODE:004516B8      mov     fs, di
CODE:004516BB      popf
CODE:004516BC      popa
CODE:004516BD      sti
CODE:004516BE      retf
```

If you observe this code you may see how it is acting like callgate dispatcher and will execute procedure that is stored in variable callgate_current, there are only 3 calls to callgate dispatcher, also cli and sti are used to disable task switching while we are in callgate.

Before we enter into process hiding techniqe we are going to see also how it gets needed offsets for EPROCESS struct and also some addresses from maped ntoskrnl.exe here:



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```
CODE:00451D50 win_xp:
CODE:00451D50          mov     ds:ActiveProcessLinks, 88h
CODE:00451D5A          mov     ds:UniqueProcessId, 84h
CODE:00451D64          mov     ds:ImageFileName, 174h
CODE:00451D6E          mov     ds:InheritedFromUniqueProcessId, 14Ch
CODE:00451D78          mov     ds:SeAuditProcessCreationInfo, 1F4h
CODE:00451D82
CODE:00451D82 loc_451D82:
CODE:00451D82          mov     eax, offset aNtoskrnl_exe ; "ntoskrnl.exe"
CODE:00451D87          call   getntoskrnlbase
CODE:00451D8C          mov     ds:real_ntoskrnl_base, eax
CODE:00451D91          push   1 ; dwFlags
CODE:00451D93          push   0 ; hFile
CODE:00451D95          push   offset aNtoskrnl_exe ; "ntoskrnl.exe"
CODE:00451D9A          call   LoadLibraryExA_0
CODE:00451D9F          mov     ds:mapped_ntoskrnl_base, eax
CODE:00451DA4          mov     eax, offset aMmgetphysicala ; "MmGetPhysicalAddress"
CODE:00451DA9          call   find_exports
CODE:00451DAE          mov     ds:MmGetPhysicalAddress, eax
CODE:00451DB3          mov     eax, offset aMmisaddressval ; "MmIsAddressValid"
CODE:00451DB8          call   find_exports
CODE:00451DBD          mov     ds:MmIsAddressValid, eax
CODE:00451DC2          mov     eax, offset aIogetcurrentpr ; "IoGetCurrentProcess"
CODE:00451DC7          call   find_exports
CODE:00451DCC          mov     ds:IoGetCurrentProcess, eax
CODE:00451DD1          mov     eax, offset aKe386setioacce ; "Ke386SetIoAccessMap"
CODE:00451DD6          call   find_exports
CODE:00451DDB          mov     ds:Ke386SetIoAccessMap, eax
CODE:00451DE0          mov     eax, offset aKe386queryioac ; "Ke386QueryIoAccessMap"
CODE:00451DE5          call   find_exports
CODE:00451DEA          mov     ds:Ke386QueryIoAccessMap, eax
CODE:00451DEF          mov     eax, offset aKe386iosetacce ; "Ke386IoSetAccessProcess"
CODE:00451DF4          call   find_exports
CODE:00451DF9          mov     ds:Ke386IoSetAccessProcess, eax
CODE:00451DFE          mov     eax, offset aExallocatepool ; "ExAllocatePool"
CODE:00451E03          call   find_exports
CODE:00451E08          mov     ds:ExAllocatePool, eax
CODE:00451E0D          mov     eax, offset aExfreepool ; "ExFreePool"
CODE:00451E12          call   find_exports
CODE:00451E17          mov     ds:ExFreePool, eax
CODE:00451E1C          mov     ds:dword_455C14, esi
CODE:00451E22          mov     eax, ds:dword_455C14
CODE:00451E27          sub     eax, 1
```

Here you may see how it stores offsets to important part of EPROCESS depending on it if it is win2k or XP, also it will locate address of ntoskrnl.exe in memory, and addresses of important exports of ntoskrnl.exe, non of them is rebased according to base of real ntoskrnl.exe because only one of them is used and that is IoGetCurrentProcess but this one only is only retrieving EPROCESS of current process so it doesn't require rebaseing. I mean it can be called from ntoskrnl.exe mapped in crackme.

Oki now lets examine calls to callgate, there is only 3 of them, and those are called from here:

```
CODE:00451A48 locate_process proc near
CODE:00451A48
CODE:00451A48 var_8 = dword ptr -8
CODE:00451A48 var_4 = dword ptr -4
CODE:00451A48
CODE:00451A48          add     esp, 0FFFFFFF8h
CODE:00451A4B          mov     dword ptr [esp], offset EPROCESS
CODE:00451A52          mov     edx, esp
CODE:00451A54          mov     eax, offset IoGetCurrentProcess_callgate
CODE:00451A59          call   call_callgate
CODE:00451A5E          mov     eax, [esp+4]
CODE:00451A62          pop     ecx
CODE:00451A63          pop     edx
CODE:00451A64          retn
CODE:00451A64 locate_process endp
```



First we receive EPROCESS of, I suppose of SMSS.exe, only parentless process is SMSS.exe created by ntoskrnl.exe, for more refer to Barnaby Jack's article about "Step Into The Ring0" [2].

```
CODE:00451A24 IoGetCurrentProcess_callgate:
CODE:00451A24         mov     ebx, eax           ; stack location with args in ring3
CODE:00451A26         call   ds:IoGetCurrentProcess
CODE:00451A2C         mov     edx, [ebx]        ; Arguments buffer
CODE:00451A2E         mov     esi, [edx+eprocess_struct.ep_activeprocesslinks]
CODE:00451A31         mov     edi, [edx+eprocess_struct.ep_InheritedFromUniqueProcessId]
CODE:00451A34 __cycle:
CODE:00451A34         mov     ecx, [edi+eax]
CODE:00451A37         test   ecx, ecx
CODE:00451A39         jz     short __exit_callgate_0
CODE:00451A3B         mov     eax, [esi+eax]    ; next EPROCESS in list entry
CODE:00451A3E         sub     eax, esi          ; negative offset to start of EPROCESS
CODE:00451A40         jmp    short __cycle
CODE:00451A42 __exit_callgate_0:
CODE:00451A42         mov     [ebx+4], eax     ; save EPROCESS
CODE:00451A45         retn
```

Next one is for locating EPROCESS of given PID, but this time PID is current process:

```
CODE:00451FEC findprocess_callgate:
CODE:00451FEC         mov     ebx, [eax]        ; EPROCESS struct
CODE:00451FEE         mov     ecx, [eax+4]     ; Current process ID
CODE:00451FF1         push   eax               ; save arguments list
CODE:00451FF2         mov     eax, [ebx]       ; eax = EPROCESS
CODE:00451FF4         mov     esi, [ebx+eprocess_struct.ep_activeprocesslinks]
CODE:00451FF7         mov     edi, [ebx+eprocess_struct.ep_uniqueprocessid]
CODE:00451FFA __find_my_eprocess:
CODE:00451FFA         mov     edx, [edi+eax]   ; get process ID from EPROCESS
CODE:00451FFD         cmp    edx, ecx         ; same as PID of process that we are
CODE:00451FFF         jz     short __found_my_eprocess ; looking for
CODE:00452001         mov     eax, [esi+eax]   ; ActiveProcessLinks.Flink
CODE:00452004         sub     eax, esi        ; start of EPROCESS
CODE:00452006         cmp    eax, [ebx]       ; is it same EPROCESS as from where
CODE:00452008         jz     short __exit_find_eprocess ;we have started
CODE:0045200A         jmp    short __find_my_eprocess
CODE:0045200C __found_my_eprocess:
CODE:0045200C         pop    edx
CODE:0045200D         mov    [edx+8], eax
CODE:00452010         retn
CODE:00452011 __exit_find_eprocess:
CODE:00452011         pop    edx
CODE:00452012         mov    dword ptr [edx+8], 0
CODE:00452019         retn
```

Unlinking occurs in 3rd callgate:

```
CODE:00452040 erase_process_callgate:
CODE:00452040         push   eax
CODE:00452041         mov    eax, [eax+4]     ; my EPROCESS
CODE:00452044         push   eax
CODE:00452045         call  ds:MmIsAddressValid
CODE:0045204B         test   eax, eax
CODE:0045204D         jz     short __wrong_address
CODE:0045204F         pop    eax
CODE:00452050         mov    ebx, [eax]       ; struct offset
CODE:00452052         mov    ecx, [eax+4]    ; EPROCESS
CODE:00452055         mov    esi, [ebx+eprocess_struct.ep_activeprocesslinks]
CODE:00452058         mov    edx, [esi+ecx]  ; EDX next EPROCESS
CODE:0045205B         add    esi, 4
CODE:0045205E         mov    edi, [esi+ecx]  ; ESI previous EPROCESS
CODE:00452061         mov    [edx+LIST_ENTRY.Blink], edi ; unlink EPROCESS
CODE:00452064         mov    [edi+LIST_ENTRY.Flink], edx ; unlink EPROCESS
```



```
CODE:00452066          retn
CODE:00452067
CODE:00452067  __wrong_address:
CODE:00452067          pop     eax
CODE:00452068          retn
```

All you can do now to make process vivible is to patch rotuines that open `\Device\PhysicalMemory` at this place:

```
CODE:004528D5          call   preapare_for_callgate ; patch here
CODE:004528DA          call   GetCurrentProcessId
CODE:004528DF          call   unlink_process ; and here
```

Now we are finaly done and we have removed all anti-debug from this crackme, really nice and chalanging crackme. Real master piece.

7. Conclusion

I hope that someone will benefit from this essay, I don't have much to say in this chapter except that I had really fun solving this crackme. I hope you have learnt something new by reading this solution.

8. References

- [1] Playing with Windows `/dev/(k)mem`, crazylord,
<http://www.phrack.org/archives/59/p59-0x10.txt>
- [2] Remote Windows Kernel Exploitation – Step Into the Ring 0, Barnaby Jack,
<http://research.eeye.com/html/papers/>

S verom u Boga, deroko/ARTeam

All the code provided with this tutorial is free for public use, just make a greetz to the authors and the ARTeam if you find it useful to use. Don't use these concepts for making illegal operation, all the info here reported are only meant for studying and to help having a better knowledge of application code security techniques.



9. Greetings

I wish to tank all the ARTeam members for sharing their knowledge, to 29a virus writing group for one of the best e-zines, Ms-Rem for this crackme... and of course you for reading this article.



<http://cracking.accessroot.com>